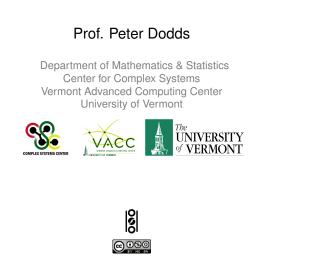
Random Networks Complex Networks, CSYS/MATH 303, Spring, 2010



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Random networks

Pure, abstract random networks:

- Consider set of all networks with N labelled nodes and *m* edges.
- Standard random network = randomly chosen network from this set.
- To be clear: each network is equally probable.
- Sometimes equiprobability is a good assumption, but it is always an assumption.
- Known as Erdős-Rényi random networks or ER graphs.

Outline Basics Structure

Random Networks

Basics

Structure

Frame 1/89

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Random Networks

Basics

Definitions

Structure

Frame 4/89

P

Clustering **Degree distributions** Configuration model Largest component

Definitions

How to build

Some visual examples

Generating Functions

Definitions **Basic Properties Giant Component Condition** Component sizes Useful results Size of the Giant Component Average Component Size References

Random networks

Some features:

Number of possible edges:

$$0 \le m \le \binom{N}{2} = \frac{N(N-1)}{2}$$

- Given *m* edges, there are $\binom{\binom{N}{2}}{m}$ different possible networks.
- Crazy factorial explosion for $1 \ll m \ll \binom{N}{2}$.
- Limit of m = 0: empty graph.
- Limit of $m = \binom{N}{2}$: complete or fully-connected graph.
- Real world: links are usually costly so real networks are almost always sparse.

Basics

Random Networks

Structure

Generating

Frame 2/89 B 990

Basics Definitions

Structure

Random Network

Generating unctions

Random networks

How to build standard random networks:

- ► Given N and m.
- Two probablistic methods (we'll see a third later on)
- 1. Connect each of the $\binom{N}{2}$ pairs with appropriate probability p.
 - Useful for theoretical work.
- 2. Take N nodes and add exactly m links by selecting edges without replacement.
 - Algorithm: Randomly choose a pair of nodes i and j, $i \neq j$, and connect if unconnected; repeat until all m edges are allocated.
 - Best for adding relatively small numbers of links (most cases).
 - ▶ 1 and 2 are effectively equivalent for large N.



Random networks A few more things:

► For method 1, # links is probablistic:

$$\langle m \rangle = p \binom{N}{2} = p \frac{1}{2} N(N-1)$$

So the expected or average degree is

$$\langle k \rangle = \frac{2 \langle m \rangle}{N}$$
$$= \frac{2}{N} p \frac{1}{2} N(N-1) = \frac{2}{N} p \frac{1}{2} M(N-1) = p(N-1)$$

- Which is what it should be...
- If we keep $\langle k \rangle$ constant then $p \propto 1/N \rightarrow 0$ as $N \to \infty$.

Random Networks

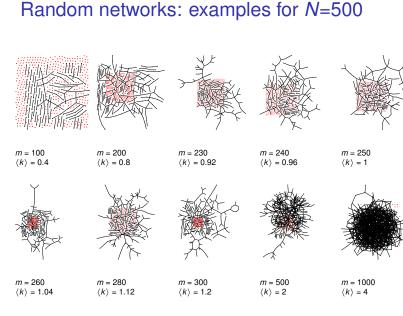
Some visual example:

Generating

Functions

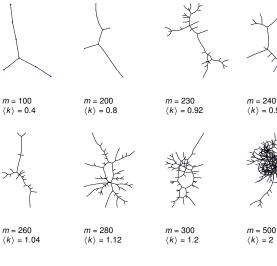
Basics

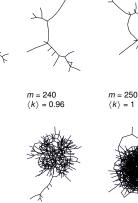
Random Networks





Random networks: largest components





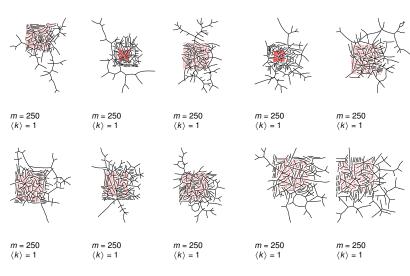
 $\langle k \rangle = 2$



m = 1000 $\langle k \rangle = 4$

Frame 22/89

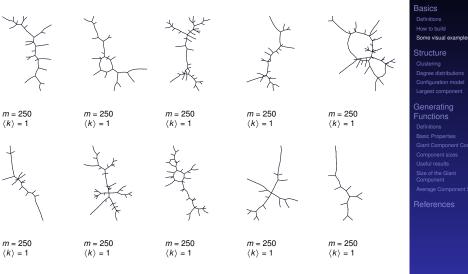
Random networks: examples for N=500





Random Networks

Random networks: largest components



Frame 24/89

Random Networks

Random networks

Clustering:

- For method 1, what is the clustering coefficient for a finite network?
- Consider triangle/triple clustering coefficient (Newman^[1]):

$$C_2 = \frac{3 \times \# \text{triangles}}{\# \text{triples}}$$

- Recall: C_2 = probability that two nodes are connected given they have a friend in common.
- ► For standard random networks, we have simply that

$$C_2 = p.$$

Random Networks
Basics
Structure
Clustering
Generating
Functions
Size of the Giant Component
Frame 26/89
- 1 - 1

Random networks

Clustering:

- ▶ So for large random networks ($N \rightarrow \infty$), clustering drops to zero.
- Key structural feature of random networks is that they locally look like branching networks (no loops).

Structure Clustering

Random Network

Basics

Generating

Random networks

Degree distribution:

- Recall p_k = probability that a randomly selected node has degree k.
- Consider method 1 for constructing random networks: each possible link is realized with probability p.
- Now consider one node: there are 'N 1 choose k' ways the node can be connected to k of the other N 1 nodes.
- Each connection occurs with probability p, each non-connection with probability (1 – p).
- Therefore have a binomial distribution:

$$P(k; p, N) = \binom{N-1}{k} p^k (1-p)^{N-1-k}$$

Limiting form of P(k; p, N):

Substitute
$$p = \frac{\langle k \rangle}{N-1}$$
 into $P(k; p, N)$ and hold k fixed:

$$P(k; p, N) = {\binom{N-1}{k}} \left(\frac{\langle k \rangle}{N-1}\right)^k \left(1 - \frac{\langle k \rangle}{N-1}\right)^{N-1-k}$$

$$= \frac{(N-1)!}{k!(N-1-k)!} \frac{\langle k \rangle^k}{(N-1)^k} \left(1 - \frac{\langle k \rangle}{N-1}\right)^{N-1-k}$$

$$= \frac{(N-1)(N-2)\cdots(N-k)}{k!} \frac{\langle k \rangle^k}{(N-1)^k} \left(1 - \frac{\langle k \rangle}{N-1}\right)^{N-1-k}$$

$$\simeq \frac{N^k(1 - \frac{\sqrt{N}}{N})\cdots(1 - \frac{\sqrt{N}}{N})}{k!N^k} \frac{\langle k \rangle^k}{(1 - \frac{\sqrt{N}}{N})^k} \left(1 - \frac{\langle k \rangle}{N-1}\right)^{N-1-k}$$



Random Networks

Basics

Structure

Degree distributions

Random networks

Limiting form of P(k; p, N):

- Our degree distribution: $P(k; p, N) = \binom{N-1}{k} p^k (1-p)^{N-1-k}.$
- What happens as $N \to \infty$?
- We must end up with the normal distribution right?
- If p is fixed, then we would end up with a Gaussian with average degree ⟨k⟩ ≃ pN → ∞.
- But we want to keep $\langle k \rangle$ fixed...
- ► So examine limit of P(k; p, N) when $p \to 0$ and $N \to \infty$ with $\langle k \rangle = p(N-1) = \text{constant.}$

Basics Definitions How to build Some visual examples **Structure** Clustering **Degree distributions** Configuration model Largest component **Generating Functions** Definitions Basic Properties Giant Component Condition Component Sizes Useful results Size of the Giant Component Size of the Giant Component Average Component Size

Random Networks

Frame 30/89 日 のへへ

Random Network

Limiting form of P(k; p, N):

► We are now here:

$$P(k; p, N) \simeq rac{\langle k
angle^k}{k!} \left(1 - rac{\langle k
angle}{N-1}
ight)^{N-1-k}$$

Now use the excellent result:

$$\lim_{n\to\infty}\left(1+\frac{x}{n}\right)^n=e^x$$

(Use l'Hôpital's rule to prove.)

• Identifying n = N - 1 and $x = -\langle k \rangle$:

$$P(k;\langle k\rangle) \simeq \frac{\langle k\rangle^k}{k!} e^{-\langle k\rangle} \left(1 - \frac{\langle k\rangle}{N-1}\right)^{-k} \rightarrow \frac{\langle k\rangle^k}{k!} e^{-\langle k\rangle}$$

▶ This is a Poisson distribution (\boxplus) with mean $\langle k \rangle$.

Basics Definitions How to build Some visual exampl Structure Olustering Degree distributions Configuration model Largest component Generating Functions Definitions

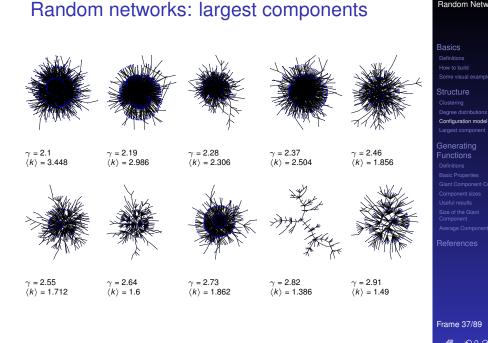
Frame <u>32/89</u>

General random networks

- So... standard random networks have a Poisson degree distribution
- Generalize to arbitrary degree distribution P_k .
- Also known as the configuration model^[1].
- Can generalize construction method from ER random networks.
- Assign each node a weight w from some distribution P_w and form links with probability

P(link between *i* and *j*) $\propto w_i w_i$.

- But we'll be more interested in
 - 1. Randomly wiring up (and rewiring) already existing nodes with fixed degrees.
 - 2. Examining mechanisms that lead to networks with certain degree distributions.

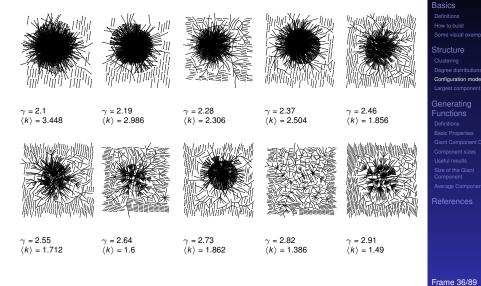


Basics Configuration model $\gamma = 2.1$ $\langle k \rangle = 3.448$ $\gamma = 2.55$ $\langle k \rangle = 1.712$ Frame 34/89

Random Networks

Random Networks

Random networks: examples for N=1000



Poisson basics:

Normalization: we must have

$$\sum_{k=0}^{\infty} P(k; \langle k \rangle) = 1$$

Checking:

$$\sum_{k=0}^{\infty} P(k; \langle k \rangle) = \sum_{k=0}^{\infty} \frac{\langle k \rangle^k}{k!} e^{-\langle k \rangle}$$
$$= e^{-\langle k \rangle} \sum_{k=0}^{\infty} \frac{\langle k \rangle^k}{k!}$$
$$= e^{-\langle k \rangle} e^{\langle k \rangle} = \mathbf{1} \checkmark$$

Structure Configuration mode Generating

Basics

Random Network

Random Networks

Frame 38/89

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Poisson basics:

Mean degree: we must have

$$\langle k \rangle = \sum_{k=0}^{\infty} k P(k; \langle k \rangle).$$

Checking:

$$\sum_{k=0}^{\infty} k P(k; \langle k \rangle) = \sum_{k=0}^{\infty} k \frac{\langle k \rangle^{k}}{k!} e^{-\langle k \rangle}$$
$$= e^{-\langle k \rangle} \sum_{k=1}^{\infty} \frac{\langle k \rangle^{k}}{(k-1)!}$$
$$= \langle k \rangle e^{-\langle k \rangle} \sum_{k=1}^{\infty} \frac{\langle k \rangle^{k-1}}{(k-1)!}$$
$$\langle k \rangle e^{-\langle k \rangle} \sum_{i=0}^{\infty} \frac{\langle k \rangle^{i}}{i!} = \langle k \rangle e^{-\langle k \rangle} e^{\langle k \rangle} = \langle k \rangle \checkmark$$

We'll get to a better way of doing this...

=

The edge-degree distribution:

- The degree distribution P_k is fundamental for our description of many complex networks
- Again: P_k is the degree of randomly chosen node.
- A second very important distribution arises from choosing randomly on edges rather than on nodes.
- Define Q_k to be the probability the node at a random end of a randomly chosen edge has degree k.
- Now choosing nodes based on their degree (i.e., size):
- Normalized form:

$$Q_k = \frac{kP_k}{\sum_{k'=0}^{\infty} k'P_{k'}} = \frac{kP_k}{\langle k \rangle}.$$

 $Q_k \propto k P_k$

Random Networks
Basics
Definitions
How to build
Some visual examples
Structure
Clustering
Degree distributions
Configuration model
Largest component
Generating
Functions
Definitions
Basic Properties
Giant Component Condition
Component sizes
Useful results
Size of the Giant Component
Average Component Size
References
Frame 39/89
- ଅ - ୬.୧.୧

Random Networks

Basics

Configuration mode

Generating

Frame 41/89

Poisson basics:

- The variance of degree distributions for random networks turns out to be very important.
- Use calculation similar to one for finding (k) to find the second moment:

$$\langle k^2 \rangle = \langle k \rangle^2 + \langle k \rangle$$

Variance is then

$$\sigma^{2} = \langle \mathbf{k}^{2} \rangle - \langle \mathbf{k} \rangle^{2} = \langle \mathbf{k} \rangle^{2} + \langle \mathbf{k} \rangle - \langle \mathbf{k} \rangle^{2} = \langle \mathbf{k} \rangle.$$

- So standard deviation σ is equal to $\sqrt{\langle k \rangle}$.
- Note: This is a special property of Poisson distribution and can trip us up...

Random Network

Basics

Configuration mod

Generating

Random Network

Basics

Structure

Configuration mod Largest componer Generating

The edge-degree distribution:

- For random networks, Q_k is also the probability that a friend (neighbor) of a random node has k friends.
- Useful variant on Q_k :

 R_k = probability that a friend of a random node has *k* other friends.

$$R_k = rac{(k+1)P_{k+1}}{\sum_{k'=0}(k'+1)P_{k'+1}} = rac{(k+1)P_{k+1}}{\langle k
angle}$$

- Equivalent to friend having degree k + 1.
- Natural question: what's the expected number of other friends that one friend has?

Frame 42/89

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The edge-degree distribution:

Given R_k is the probability that a friend has k other friends, then the average number of friends' other friends is

$$\langle k \rangle_R = \sum_{k=0}^{\infty} kR_k = \sum_{k=0}^{\infty} k \frac{(k+1)P_{k+1}}{\langle k \rangle}$$
$$= \frac{1}{\langle k \rangle} \sum_{k=1}^{\infty} k(k+1)P_{k+1}$$
$$= \frac{1}{\langle k \rangle} \sum_{k=1}^{\infty} \left((k+1)^2 - (k+1) \right) P_{k+1}$$

(where we have sneakily matched up indices)

$$= \frac{1}{\langle k \rangle} \sum_{j=0}^{\infty} (j^2 - j) P_j \quad \text{(using j = k+1)}$$
$$= \frac{1}{\langle k \rangle} \left(\langle k^2 \rangle - \langle k \rangle \right)$$

Two reasons why this matters

Reason #1:

Average # friends of friends per node is

$$\langle k_2 \rangle = \langle k \rangle \times \langle k \rangle_R = \langle k \rangle \frac{1}{\langle k \rangle} \left(\langle k^2 \rangle - \langle k \rangle \right) = \langle k^2 \rangle - \langle k \rangle.$$

- Key: Average depends on the 1st and 2nd moments of P_k and not just the 1st moment.
- Three peculiarities:
 - 1. We might guess $\langle k_2 \rangle = \langle k \rangle (\langle k \rangle 1)$ but it's actually $\langle k(k-1) \rangle$.
 - 2. If P_k has a large second moment, then $\langle k_2 \rangle$ will be big.
 - (e.g., in the case of a power-law distribution)
 - 3. Your friends are different to you...

Handon Networks
Basics
Definitions
How to build
Some visual examples
Structure
Clustering
Degree distributions
Configuration model
Largest component
Generating
Functions
Definitions
Basic Properties
Giant Component Condition
Component sizes
Useful results
Size of the Giant
Component
Average Component Size
References
11010101000
Frame 43/89
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Random Networks

Basics

Definitions

Structure

Configuration mode

Generating Functions

Frame 4<u>5/89</u>

P

The edge-degree distribution:

- Note: our result, ⟨k⟩_R = 1/⟨k⟩ (⟨k²⟩ ⟨k⟩), is true for all random networks, independent of degree distribution.
- ▶ For standard random networks, recall

$$\langle k^2 \rangle = \langle k \rangle^2 + \langle k \rangle.$$

Therefore:

$$\left\langle k \right\rangle_{R} = \frac{1}{\left\langle k \right\rangle} \left(\left\langle k \right\rangle^{2} + \left\langle k \right\rangle - \left\langle k \right\rangle \right) = \left\langle k \right\rangle$$

- Again, neatness of results is a special property of the Poisson distribution.
- ► So friends on average have $\langle k \rangle$ other friends, and $\langle k \rangle$ + 1 total friends...

Random Networks

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Random Network

Basics

Frame 44/89

Two reasons why this matters

More on peculiarity #3:

- A node's average # of friends: (k)
- Friend's average # of friends: $\frac{\langle k^2 \rangle}{\langle k \rangle}$
- Comparison:

$$\frac{\langle k^2 \rangle}{\langle k \rangle} = \langle k \rangle \frac{\langle k^2 \rangle}{\langle k \rangle^2} = \langle k \rangle \frac{\sigma^2 + \langle k \rangle^2}{\langle k \rangle^2} = \langle k \rangle \left(1 + \frac{\sigma^2}{\langle k \rangle^2} \right) \ge$$

- So only if everyone has the same degree (variance= σ² = 0) can a node be the same as its friends.
- Intuition: for random networks, the more connected a node, the more likely it is to be chosen as a friend.

Some visual example Structure Clustering Degree distributions Configuration model Largest component Generating Functions Definitions

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Frame 46/89

Two reasons why this matters

(Big) Reason #2:

- \$\langle k \rangle_R\$ is key to understanding how well random networks are connected together.
- e.g., we'd like to know what's the size of the largest component within a network.
- As N → ∞, does our network have a giant component?
- Defn: Component = connected subnetwork of nodes such that ∃ path between each pair of nodes in the subnetwork, and no node outside of the subnetwork is connected to it.
- Defn: Giant component = component that comprises a non-zero fraction of a network as N → ∞.
- Note: Component = Cluster

Giant component

Standard random networks:

- Recall $\langle k^2 \rangle = \langle k \rangle^2 + \langle k \rangle$.
- Condition for giant component:

$$\langle k \rangle_{R} = \frac{\langle k^{2} \rangle - \langle k \rangle}{\langle k \rangle} = \frac{\langle k \rangle^{2} + \langle k \rangle - \langle k \rangle}{\langle k \rangle} = \langle k \rangle$$

- Therefore when (k) > 1, standard random networks have a giant component.
- When $\langle k \rangle < 1$, all components are finite.
- ► Fine example of a continuous phase transition (⊞).
- We say $\langle k \rangle = 1$ marks the critical point of the system.



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Random Networks

Basics

Definitions

Structure

Largest component

Generating Functions

Frame 50/89

Structure of random networks

Giant component:

- A giant component exists if when we follow a random edge, we are likely to hit a node with at least 1 other outgoing edge.
- Equivalently, expect exponential growth in node number as we move out from a random node.
- All of this is the same as requiring $\langle k \rangle_R > 1$.
- Giant component condition (or percolation condition):

$$\langle k \rangle_R = rac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle} > 1$$

- Again, see that the second moment is an essential part of the story.
- Equivalent statement: $\langle k^2 \rangle > 2 \langle k \rangle$

Frame 49/89 奇 のへで

Random Network

Random Network

Basics

Structure

Largest componen

Generating

unctions

Giant component

Random networks with skewed P_k :

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• e.g, if
$$P_k = ck^{-\gamma}$$
 with 2 < γ < 3 then

$$\langle k^2
angle = c \sum_{k=0}^{\infty} k^2 k^{-\gamma}$$

$$\sim \int_{x=0}^{\infty} x^{2-\gamma} \mathrm{d}x$$
$$x^{3-\gamma}\Big|_{x=0}^{\infty} = \infty \quad (>\langle k \rangle).$$

- So giant component always exists for these kinds of networks.
- Cutoff scaling is k⁻³: if γ > 3 then we have to look harder at ⟨k⟩_B.

Basics Definitions How to build Some visual example Structure Clustering Degree distributions Configuration model Largest component Generating Functions Definitions

> Biant Component Co Component sizes Useful results Bize of the Giant Component

> > erences

Giant component

And how big is the largest component?

- Define S_1 as the size of the largest component.
- Consider an infinite ER random network with average degree (k).
- Let's find S_1 with a back-of-the-envelope argument.
- Define δ as the probability that a randomly chosen node does not belong to the largest component.
- Simple connection: $\delta = 1 S_1$.
- Dirty trick: If a randomly chosen node is not part of the largest component, then none of its neighbors are.
- So

$$\delta = \sum_{k=0}^{\infty} P_k \delta^k$$

Substitute in Poisson distribution...

Giant component

- We can figure out some limits and details for $S_1 = 1 e^{-\langle k \rangle S_1}$.
- First, we can write $\langle k \rangle$ in terms of S_1 :

$$\langle k \rangle = rac{1}{S_1} \ln rac{1}{1-S_1}$$

- ▶ As $\langle k \rangle \rightarrow 0$, $S_1 \rightarrow 0$.
- As $\langle k \rangle \rightarrow \infty$, $S_1 \rightarrow 1$.
- Notice that at $\langle k \rangle = 1$, the critical point, $S_1 = 0$.
- Only solvable for S > 0 when $\langle k \rangle > 1$.
- ▶ Really a transcritical bifurcation^[2].

Giant component

Carrying on:

$$\delta = \sum_{k=0}^{\infty} P_k \delta^k = \sum_{k=0}^{\infty} \frac{\langle k \rangle^k}{k!} e^{-\langle k \rangle} \delta^k$$
$$= e^{-\langle k \rangle} \sum_{k=0}^{\infty} \frac{(\langle k \rangle \delta)^k}{k!}$$
$$= e^{-\langle k \rangle} e^{\langle k \rangle \delta} = e^{-\langle k \rangle (1-\delta)}.$$

▶ Now substitute in $\delta = 1 - S_1$ and rearrange to obtain:

$$S_1 = 1 - e^{-\langle k \rangle S_1}$$

Frame 53/89 日 つくへ

Random Networks

Basics

Structure

Largest component Generating

Random Networks

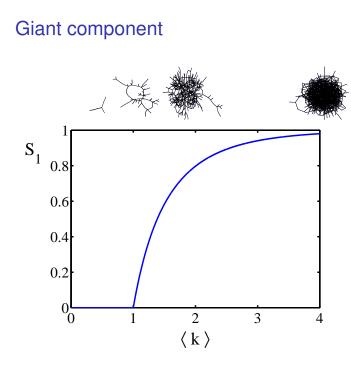
Random Networks

Basics

Structure

Largest component

Frame 52/89



Random Network

Basics Definitions How to build Some visual example Structure

Degree distributions Configuration model Largest component Generating

Functions Definitions Basic Properties Giant Component Ci Component sizes Useful results

ferences

Frame 55/89

Giant component

Turns out we were lucky ...

- Our dirty trick only works for ER random networks.
- The problem: We assumed that neighbors have the same probability δ of belonging to the largest component.
- But we know our friends are different from us...
- Works for ER random networks because $\langle k \rangle = \langle k \rangle_R$.
- We need a separate probability δ' for the chance that a node at the end of a random edge is part of the largest component.
- We can do this but we need to enhance our toolkit with Generatingfunctionology...^[3]
- (Well, not really but it's fun and we get all sorts of other things...)

Simple example

Rolling dice:

$$F^{(\Box)}(x) = \sum_{k=1}^{6} p_k x^k = \frac{1}{6} (x + x^2 + x^3 + x^4 + x^5 + x^6)$$

We'll come back to this simple example as we derive various delicious properties of generating functions.



Random Networks

Basics

Definitions

Structure

Generating Functions

Frame 59/89

P

Definitions

Generating functions

- Idea: Given a sequence a₀, a₁, a₂,..., associate each element with a distinct function or other mathematical object.
- Well-chosen functions allow us to manipulate sequences and retrieve sequence elements.

Definition:

• The generating function (g.f.) for a sequence $\{a_n\}$ is

$$F(x) = \sum_{n=0}^{\infty} a_n x^n.$$

- ► Roughly: transforms a vector in R[∞] into a function defined on R¹.
- Related to Fourier, Laplace, Mellin, ...

Random Network

Random Networks

Basics

Structure

Generating

Definitions

Example

Take a degree distribution with exponential decay:

 $P_k = ce^{-\lambda k}$

where $c = 1 - e^{-\lambda}$.

> The generating function for this distribution is

$$F(x) = \sum_{k=0}^{\infty} P_k x^k = \sum_{k=0}^{\infty} c e^{-\lambda k} x^k = \frac{c}{1 - x e^{-\lambda}}$$

- Notice that $F(1) = c/(1 e^{-\lambda}) = 1$.
- For probability distributions, we must always have
 F(1) = 1 since

$$F(1) = \sum_{k=0}^{\infty} P_k 1^k = \sum_{k=0}^{\infty} P_k = 1.$$

Basics
Definitions
How to build
Some visual examples
Structure
Clustering
Degree distributions
Configuration model
Largest component
Generating
Functions
Definitions
Basic Properties
Giant Component size
Size of the Giant
Component Size

Frame 60/89

ন্দ ৩৫৫

Properties of generating functions

► Average degree:

$$\langle k \rangle = \sum_{k=0}^{\infty} k P_k = \sum_{k=0}^{\infty} k P_k x^{k-1} \Big|_{x=1}$$

= $\frac{d}{dx} F(x) \Big|_{x=1} = F'(1)$

- In general, many calculations become simple, if a little abstract.
- ► For our exponential example:

$$F'(x)=\frac{(1-e^{-\lambda})e^{-\lambda}}{(1-xe^{-\lambda})^2}.$$

So:

$$\langle k
angle = F'(1) = rac{e^{-\lambda}}{(1-e^{-\lambda})}$$

Edge-degree distribution

Recall our condition for a giant component:

$$\langle k \rangle_R = \frac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle} > 1.$$

- Let's r
 r
 express our condition in terms of generating functions.
- We first need the g.f. for R_k .
- ► We'll now use this notation:

 $F_P(x)$ is the g.f. for P_k . $F_B(x)$ is the g.f. for R_k .

Condition in terms of g.f. is:

$$\langle k \rangle_R = F'_R(1) > 1$$

• Now find how F_R is related to F_P ...

Random Networks
Basics
Structure
Generating
Functions
Basic Properties
Size of the Giant Component
References
Frame 62/89
- A - A & A

Properties of generating functions

Useful pieces for probability distributions:

- Normalization:
- First moment:

$$\langle k \rangle = F'(1)$$

F(1) = 1

Higher moments:

$$\langle k^n \rangle = \left(x \frac{\mathrm{d}}{\mathrm{d}x} \right)^n F(x) \Big|_{x=1}$$

kth element of sequence (general):

 $P_k = \frac{1}{k!} \frac{\mathrm{d}^k}{\mathrm{d}x^k} F(x) \bigg|_{x=0}$

How to build Some visual examples Structure Clustering Degree distributions Configuration model Largest component Generating Functions Basic Properties Giant Component Conditio Component sizes Useful results Size of the Cliant Component Size References Frame 63/89

Random Network

Basics

Structure

Generating

Giant Component Condi

Random Networks

Basics

Random Networks Basics Definitions How to build Some visual examples Clustering Degree distributions Configuration model Largest component Configuration model Largest component Configuration model Largest component Component Sizes Useful results Size of the Glant Component Sizes Useful results Size of the Glant Component Sizes Component Size Compone

Edge-degree distribution

We have

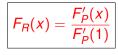
$$F_R(x) = \sum_{k=0}^{\infty} \mathbf{R}_k x^k = \sum_{k=0}^{\infty} \frac{(k+1)\mathbf{P}_{k+1}}{\langle k \rangle} x^k.$$

Shift index to j = k + 1 and pull out $\frac{1}{\langle k \rangle}$:

$$F_{R}(x) = \frac{1}{\langle k \rangle} \sum_{j=1}^{\infty} j P_{j} x^{j-1} = \frac{1}{\langle k \rangle} \sum_{j=1}^{\infty} P_{j} \frac{\mathrm{d}}{\mathrm{d}x} x^{j}$$

$$=\frac{1}{\langle k\rangle}\frac{\mathrm{d}}{\mathrm{d}x}\sum_{j=1}^{\infty}P_{j}x^{j}=\frac{1}{\langle k\rangle}\frac{\mathrm{d}}{\mathrm{d}x}\left(F_{P}(x)-P_{0}\right)=\frac{1}{\langle k\rangle}F_{P}'(x)$$

Finally, since $\langle k \rangle = F'_P(1)$,



Edge-degree distribution

- Recall giant component condition is $\langle k \rangle_R = F'_R(1) > 1.$
- Since we have $F_B(x) = F'_P(x)/F'_P(1)$,

$$F'_{R}(x) = rac{F''_{P}(x)}{F'_{P}(1).}$$

• Setting x = 1, our condition becomes



Size distributions

G.f.'s for component size distributions:

$$F_{\pi}(x) = \sum_{n=0}^{\infty} \pi_n x^n$$
 and $F_{\rho}(x) = \sum_{n=0}^{\infty} \rho_n x^n$

The largest component:

Subtle key: $F_{\pi}(1)$ is the probability that a node belongs to a finite component.

• Therefore: $S_1 = 1 - F_{\pi}(1)$.

Our mission, which we accept:

Find the four generating functions

$$F_P, F_B, F_{\pi}$$
, and F_{ρ} .



Random Networks

Basics

Structure

Functions

Component sizes

Frame 7<u>0/8</u>9

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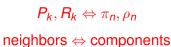
Size distributions

To figure out the size of the largest component (S_1) , we need more resolution on component sizes.

Definitions:

- $\blacktriangleright \pi_n$ = probability that a random node belongs to a finite component of size $n < \infty$.
- ρ_n = probability a random link leads to a finite subcomponent of size $n < \infty$.

Local-global connection:



Frame 69/89



Sneaky Result 1:

- Consider two random variables U and V whose values may be 0, 1, 2, ...
- Write probability distributions as U_k and V_k and g.f.'s as F_{II} and F_V .
- SR1: If a third random variable is defined as

$$W = \sum_{i=1}^{U} V^{(i)}$$
 with each $V^{(i)} \stackrel{d}{=} V$

then

$$F_W(x) = F_U(F_V(x))$$

Random Networks

Basics

Structure

Generating

Component sizes

Random Network

Basics

Useful results

Proof of SR1:

Write probability that variable W has value k as W_k .

$$W_{k} = \sum_{j=0}^{\infty} U_{j} \times \Pr(\text{sum of } j \text{ draws of variable } V = k)$$

$$= \sum_{j=0}^{\infty} U_{j} \sum_{\substack{\{i_{1}, i_{2}, \dots, i_{j}\} \mid \\ i_{1}+i_{2}+\dots+i_{j}=k}} V_{i_{1}} V_{i_{2}} \cdots V_{i_{j}}$$

$$\therefore F_{W}(x) = \sum_{k=0}^{\infty} W_{k} x^{k} = \sum_{k=0}^{\infty} \sum_{j=0}^{\infty} U_{j} \sum_{\substack{\{i_{1}, i_{2}, \dots, i_{j}\} \mid \\ i_{1}+i_{2}+\dots+i_{j}=k}} V_{i_{1}} V_{i_{2}} \sum_{j=0}^{\infty} V_{j_{1}} x^{i_{1}} V_{i_{2}} x^{i_{2}} \cdots V_{j_{j}} x^{i_{j}}$$

$$= \sum_{j=0}^{\infty} U_{j} \sum_{k=0}^{\infty} \sum_{\substack{\{i_{1}, i_{2}, \dots, i_{j}\} \mid \\ i_{1}+i_{2}+\dots+i_{j}=k}} V_{i_{1}} x^{i_{1}} V_{i_{2}} x^{i_{2}} \cdots V_{i_{j}} x^{i_{j}}}$$

Useful results we'll need for g.f.'s

Sneaky Result 2:

- Start with a random variable U with distribution U_k (k = 0, 1, 2, ...)
- SR2: If a second random variable is defined as

$$V = U + 1$$
 then $F_V(x) = xF_U(x)$

• Reason: $V_k = U_{k-1}$ for $k \ge 1$ and $V_0 = 0$.

$$\therefore F_V(x) = \sum_{k=0}^{\infty} V_k x^k = \sum_{k=1}^{\infty} U_{k-1} x^k$$
$$= x \sum_{j=0}^{\infty} U_j x^j = x F_U(x) \cdot \checkmark$$

Random Network
Basics
Structure
Clustering
Degree distributions
Generating
Functions
Useful results
Size of the Giant Component
Frame 75/89
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Random Networks

Basics

Proof of SR1:

With some concentration, observe:

Useful results we'll need for g.f.'s

Generalization of SR2:

▶ (1) If V = U + i then

▶ (2) If V = U - i then

$$F_{W}(x) = \sum_{j=0}^{\infty} U_{j} \sum_{k=0}^{\infty} \sum_{\substack{\{i_{1},i_{2},\dots,i_{j}\} \mid \\ i_{1}+i_{2}+\dots+i_{j}=k}} V_{i_{1}} x^{i_{1}} V_{i_{2}} x^{i_{2}} \dots V_{i_{j}} x^{i_{j}}} \\ \underbrace{x^{k} \text{ piece of } \left(\sum_{i'=0}^{\infty} V_{i'} x^{i'}\right)^{j}}_{\left(\sum_{i'=0}^{\infty} V_{i'} x^{i'}\right)^{j}} = (F_{V}(x))^{j}} \\ = \sum_{j=0}^{\infty} U_{j} (F_{V}(x))^{j} \\ = F_{U} (F_{V}(x)) \checkmark$$

 $F_V(x) = x^i F_U(x).$

 $F_V(x) = x^{-i} F_U(x)$

 $= x^{-i} \sum_{k=0}^{\infty} U_k x^k$

Frame 74/89

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Random Network

Basics

Structure

Generating

Useful results

Frame 76/89

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Connecting generating functions

- Goal: figure out forms of the component generating functions, *F_π* and *F_ρ*.
- π_n = probability that a random node belongs to a finite component of size n

 $= \sum_{k=0}^{\infty} P_k \times \Pr\left(\begin{array}{c} \text{sum of sizes of subcomponents} \\ \text{at end of } k \text{ random links} = n-1 \end{array}\right)$

Therefore: $F_{\pi}(x) = \underbrace{x}_{\text{SR2}} \underbrace{F_{P}(F_{\rho}(x))}_{\text{SR2}}$

Extra factor of x accounts for random node itself.

Connecting generating functions

We now have two functional equations connecting our generating functions:

 $F_{\pi}(x) = xF_{P}(F_{\rho}(x))$ and $F_{\rho}(x) = xF_{R}(F_{\rho}(x))$

- Taking stock: We know $F_P(x)$ and $F_R(x) = F'_P(x)/F'_P(1)$.
- We first untangle the second equation to find F_{ρ}
- We can do this because it only involves F_{ρ} and F_{R} .
- The first equation then immediately gives us F_π in terms of F_ρ and F_R.



Random Networks

Basics

Structure

Generating

Size of the Giant

Frame 78/89

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Functions

Connecting generating functions

- *ρ_n* = probability that a random link leads to a finite subcomponent of size *n*.
- Invoke one step of recursion: ρ_n = probability that in following a random edge, the outgoing edges of the node reached lead to finite subcomponents of combined size n 1,

$$= \sum_{k=0}^{\infty} R_k \times \Pr\left(\begin{array}{c} \text{sum of sizes of subcomponents} \\ \text{at end of } k \text{ random links} = n-1 \end{array}\right)$$

Therefore:
$$F_{\rho}(x) = \underbrace{x}_{\text{SR2}} \underbrace{F_{R}(F_{\rho}(x))}_{\text{SR1}}$$

 Again, extra factor of x accounts for random node itself. Frame 79/89

Random Network

Basics

Random Networks

Basics

Structure

Generating

Size of the Gian

unctions

Component sizes

Remembering vaguely what we are doing:

Finding F_{π} to obtain the fractional size of the largest component $S_1 = 1 - F_{\pi}(1)$.

Set x = 1 in our two equations:

 $F_{\pi}(1) = F_{P}(F_{\rho}(1))$ and $F_{\rho}(1) = F_{R}(F_{\rho}(1))$

- Solve second equation numerically for $F_{\rho}(1)$.
- Plug $F_{\rho}(1)$ into first equation to obtain $F_{\pi}(1)$.

Frame 81/89

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Component sizes

Example: Standard random graphs.

• We can show $F_P(x) = e^{-\langle k \rangle (1-x)}$

.
$$F_R(x) = F'_P(x)/F'_P(1) = e^{-\langle k \rangle (1-x)}/e^{-\langle k \rangle (1-x')}|_{x'=1}$$

 $= e^{-\langle k \rangle (1-x)} = F_P(x)$...aha!

- RHS's of our two equations are the same.
- So $F_{\pi}(x) = F_{\rho}(x) = xF_R(F_{\rho}(x)) = xF_R(F_{\pi}(x))$
- Why our dirty (but wrong) trick worked earlier...

Average component size

- Next: find average size of finite components $\langle n \rangle$.
- Using standard G.F. result: $\langle n \rangle = F'_{\pi}(1)$.
- Try to avoid finding $F_{\pi}(x)$...
- ▶ Starting from $F_{\pi}(x) = xF_{P}(F_{\rho}(x))$, we differentiate:

$$F'_{\pi}(x) = F_{\mathcal{P}}\left(F_{\rho}(x)
ight) + xF'_{
ho}(x)F'_{\mathcal{P}}\left(F_{
ho}(x)
ight)$$

• While $F_{\rho}(x) = xF_R(F_{\rho}(x))$ gives

 $F_{
ho}'(x)=F_R\left(F_{
ho}(x)
ight)+xF_{
ho}'(x)F_R'\left(F_{
ho}(x)
ight)$

- Now set x = 1 in both equations.
- We solve the second equation for F'_ρ(1) (we must already have F_ρ(1)).
- Plug $F'_{\rho}(1)$ and $F_{\rho}(1)$ into first equation to find $F'_{\pi}(1)$.

Random Networks
Basics
Structure
Degree distributions Configuration model
Largest component
Generating Functions
Definitions
Basic Properties
Size of the Giant Component
References
Frame 82/89
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Random Network

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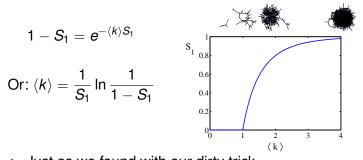
Frame 85/89

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Basics

Component sizes

- We are down to $F_{\pi}(x) = xF_{R}(F_{\pi}(x))$ and $F_{R}(x) = e^{-\langle k \rangle (1-x)}$. • $\therefore F_{\pi}(x) = xe^{-\langle k \rangle (1-F_{\pi}(x))}$
- We're first after S₁ = 1 − F_π(1) so set x = 1 and replace F_π(1) by 1 − S₁:



Just as we found with our dirty trick ...
Again, we (usually) have to resort to numerics ...

Size of the Giant Component Average Component References

Random Network

Random Network

Basics

Basics Definitions How to build Some visual examples Structure Clustering Degree distributions Configuration model Largest component Generating Functions Basic Properties Giant Component Conditi Component sizes Useful results Size of the Giant Component Sizes Hereine Component Sizes

End result: $\langle n \rangle = F'_{\pi}(1) = \frac{(1-S_1)}{1-\langle k \rangle(1-S_1)}$

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Average component size

Example: Standard random graphs.

- Use fact that $F_P = F_R$ and $F_{\pi} = F_{\rho}$.
- Two differentiated equations reduce to only one:

 $F'_{\pi}(x) = F_P\left(F_{\pi}(x)\right) + xF'_{\pi}(x)F'_P\left(F_{\pi}(x)\right)$

Rearrange:
$$F'_{\pi}(x) = \frac{F_P(F_{\pi}(x))}{1 - xF'_P(F_{\pi}(x))}$$

Simplify denominator using
$$F'_P(x) = \langle k \rangle F_P(x)$$

Frame 86

Average component size

Our result for standard random networks:

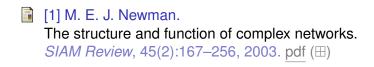
$$|n\rangle = F'_{\pi}(1) = rac{(1-S_1)}{1-\langle k
angle(1-S_1)}$$

- Recall that (k) = 1 is the critical value of average degree for standard random networks.
- Look at what happens when we increase (k) to 1 from below.
- We have $S_1 = 0$ for all $\langle k \rangle < 1$ so

$$\langle n \rangle = \frac{1}{1 - \langle k \rangle}$$

- This blows up as $\langle k \rangle \rightarrow 1$.
- Reason: we have a power law distribution of component sizes at (k) = 1.
- Typical critical point behavior....

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 Nonlinear Dynamics and Chaos.
 Addison Wesley, Reading, Massachusetts, 1994.

[3] H. S. Wilf.

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Random Networks
Basics
Definitions
How to build
Some visual examples
Structure
Clustering
Degree distributions
Configuration model
Largest component
Generating
Functions
Definitions
Basic Properties
Glant Component Condition
Component sizes
Useful results
Size of the Glant
Component Size
References

Frame 87/89

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Average component size

• Limits of $\langle k \rangle = 0$ and ∞ make sense for

$$\langle n \rangle = F'_{\pi}(1) = rac{(1-S_1)}{1-\langle k \rangle(1-S_1)}$$

- As $\langle k \rangle \rightarrow 0$, $S_1 = 0$, and $\langle n \rangle \rightarrow 1$.
- All nodes are isolated.
- As $\langle k \rangle \rightarrow \infty$, $S_1 \rightarrow 1$ and $\langle n \rangle \rightarrow 0$.
- No nodes are outside of the giant component.

Extra on largest component size:

- For $\langle k \rangle = 1$, $S_1 \sim N^{2/3}$.
- For $\langle k \rangle < 1$, $S_1 \sim \log N$.

Basics Definitions How to build Some visual examples Structure Clustering Degree distributions Configuration model Largest component Generating Functions Basic Properties Giant Component Sizes Useful results Size of the Giant Component Sizes References

Random Networks

Frame 88/89

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Random Networks