# Random Networks Complex Networks, Course 303A, Spring, 2009

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### Random networks

#### Pure, abstract random networks:

- Consider set of all networks with N labelled nodes and m edges.
- Standard random network = randomly chosen network from this set.
- To be clear: each network is equally probable.
- Sometimes equiprobability is a good assumption, but it is always an assumption.
- Known as Erdős-Rényi random networks or ER graphs.



Random Networks

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#### Outline Basics

Definitions How to build Some visual examples

#### Structure

Clustering Degree distributions Configuration model Largest component

#### **Generating Functions**

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#### Random networks

#### Some features:

Number of possible edges:

$$0 \le m \le \binom{N}{2} = \frac{N(N-1)}{2}$$

- Given *m* edges, there are 
   <sup>N</sup>
   <sup>N</sup>
- Crazy factorial explosion for  $1 \ll m \ll {N \choose 2}$ .
- Limit of m = 0: empty graph.
- Limit of  $m = \binom{N}{2}$ : complete or fully-connected graph.
- Real world: links are usually costly so real networks are almost always sparse.

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### Random networks

#### How to build standard random networks:

- ▶ Given N and m.
- Two probablistic methods (we'll see a third later on)
- 1. Connect each of the  $\binom{N}{2}$  pairs with appropriate probability *p*.
  - Useful for theoretical work.
- 2. Take *N* nodes and add exactly *m* links by selecting edges without replacement.
  - ► Algorithm: Randomly choose a pair of nodes *i* and *j*, *i* ≠ *j*, and connect if unconnected; repeat until all *m* edges are allocated.
  - Best for adding small numbers of links (most cases).
  - ▶ 1 and 2 are effectively equivalent for large *N*.



#### Random networks A few more things:

► For method 1, # links is probablistic:

$$\langle m \rangle = p \binom{N}{2} = p \frac{1}{2} N(N-1)$$

► So the expected or average degree is

$$\langle k \rangle = \frac{2 \langle m \rangle}{N}$$
$$= \frac{2}{N} p \frac{1}{2} N(N-1) = \frac{2}{N} p \frac{1}{2} M(N-1) = p(N-1)$$

- Which is what it should be...
- If we keep  $\langle k \rangle$  constant then  $p \propto 1/N \rightarrow 0$  as  $N \rightarrow \infty$ .

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Some visual example:

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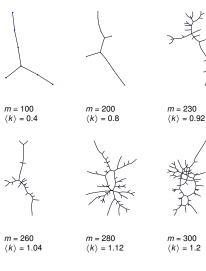
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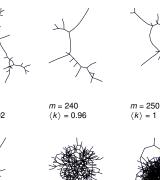
Random Networks

#### Random networks: examples for N=500 *m* = 100 *m* = 200 m = 230m = 240m = 250 $\langle k \rangle = 0.4$ $\langle k \rangle = 0.8$ $\langle k \rangle = 0.92$ $\langle k \rangle = 0.96$ $\langle k \rangle = 1$ m = 500m = 1000m = 260*m* = 280 m = 300 $\langle k \rangle = 1.2$ $\langle k \rangle = 4$ $\langle k \rangle = 1.04$ $\langle k \rangle = 1.12$ $\langle k \rangle = 2$



# Random networks: largest components





m = 500

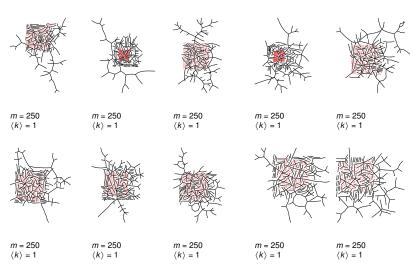
 $\langle k \rangle = 2$ 



m = 1000 $\langle k \rangle = 4$ 

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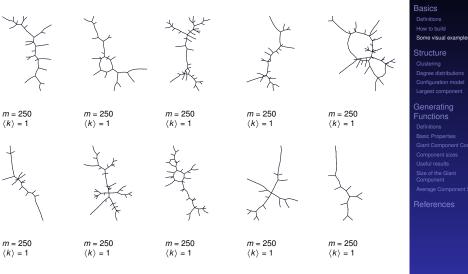
#### Random networks: examples for N=500





Random Networks

# Random networks: largest components



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Random Networks

# Random networks

#### **Clustering:**

- For method 1, what is the clustering coefficient for a finite network?
- Consider triangle/triple clustering coefficient (Newman<sup>[1]</sup>):

$$C_2 = \frac{3 \times \# \text{triangles}}{\# \text{triples}}$$

- Recall:  $C_2$  = probability that two nodes are connected given they have a friend in common.
- ► For standard random networks, we have simply that

$$C_2 = p.$$

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#### **Random networks**

#### **Clustering:**

- ▶ So for large random networks ( $N \rightarrow \infty$ ), clustering drops to zero.
- Key structural feature of random networks is that they locally look like branching networks (no loops).

Structure Clustering

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#### Random networks

#### Degree distribution:

- Recall p<sub>k</sub> = probability that a randomly selected node has degree k.
- Consider method 1 for constructing random networks: each possible link is realized with probability p.
- Now consider one node: there are 'N 1 choose k' ways the node can be connected to k of the other N 1 nodes.
- Each connection occurs with probability p, each non-connection with probability (1 – p).
- Therefore have a binomial distribution:

$$P(k; p, N) = \binom{N-1}{k} p^k (1-p)^{N-1-k}$$

# Limiting form of P(k; p, N):

Substitute 
$$p = \frac{\langle k \rangle}{N-1}$$
 into  $P(k; p, N)$  and hold  $k$  fixed:  

$$P(k; p, N) = {\binom{N-1}{k}} \left(\frac{\langle k \rangle}{N-1}\right)^k \left(1 - \frac{\langle k \rangle}{N-1}\right)^{N-1-k}$$

$$= \frac{(N-1)!}{k!(N-1-k)!} \frac{\langle k \rangle^k}{(N-1)^k} \left(1 - \frac{\langle k \rangle}{N-1}\right)^{N-1-k}$$

$$= \frac{(N-1)(N-2)\cdots(N-k)}{k!} \frac{\langle k \rangle^k}{(N-1)^k} \left(1 - \frac{\langle k \rangle}{N-1}\right)^{N-1-k}$$

$$\simeq \frac{N^k(1 - \frac{\sqrt{N}}{N})\cdots(1 - \frac{\sqrt{N}}{N})}{k!N^k} \frac{\langle k \rangle^k}{(1 - \frac{\sqrt{N}}{N})^k} \left(1 - \frac{\langle k \rangle}{N-1}\right)^{N-1-k}$$



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#### Limiting form of P(k; p, N):

- Our degree distribution:  $P(k; p, N) = \binom{N-1}{k} p^k (1-p)^{N-1-k}.$
- What happens as  $N \to \infty$ ?
- We must end up with the normal distribution right?
- If p is fixed, then we would end up with a Gaussian with average degree ⟨k⟩ ≃ pN → ∞.
- But we want to keep  $\langle k \rangle$  fixed...
- ► So examine limit of P(k; p, N) when  $p \to 0$  and  $N \to \infty$  with  $\langle k \rangle = p(N-1) = \text{constant.}$

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# Limiting form of P(k; p, N):

► We are now here:

$$P(k; p, N) \simeq rac{\langle k 
angle^k}{k!} \left(1 - rac{\langle k 
angle}{N-1}
ight)^{N-1-k}$$

Now use the excellent result:

$$\lim_{n\to\infty}\left(1+\frac{x}{n}\right)^n=e^x$$

(Use l'Hôpital's rule to prove.)

• Identifying n = N - 1 and  $x = -\langle k \rangle$ :

$$P(k;\langle k\rangle) \simeq \frac{\langle k\rangle^k}{k!} e^{-\langle k\rangle} \left(1 - \frac{\langle k\rangle}{N-1}\right)^{-k} \rightarrow \frac{\langle k\rangle^k}{k!} e^{-\langle k\rangle}$$

▶ This is a Poisson distribution ( $\boxplus$ ) with mean  $\langle k \rangle$ .

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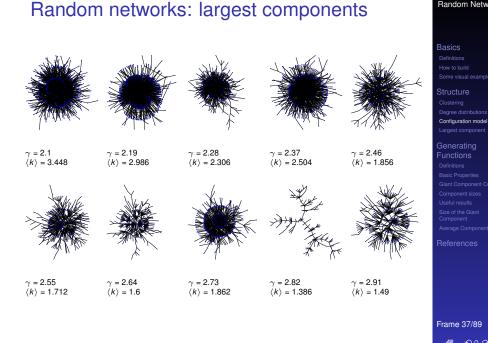
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#### General random networks

- So... standard random networks have a Poisson degree distribution
- Generalize to arbitrary degree distribution  $P_k$ .
- Also known as the configuration model<sup>[1]</sup>.
- Can generalize construction method from ER random networks.
- Assign each node a weight w from some distribution  $P_w$  and form links with probability

*P*(link between *i* and *j*)  $\propto w_i w_i$ .

- But we'll be more interested in
  - 1. Randomly wiring up (and rewiring) already existing nodes with fixed degrees.
  - 2. Examining mechanisms that lead to networks with certain degree distributions.

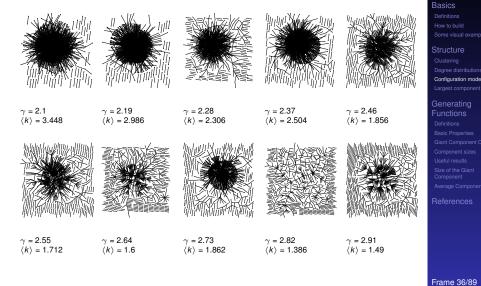


# Basics Configuration model $\gamma = 2.1$ $\langle k \rangle = 3.448$ $\gamma = 2.55$ $\langle k \rangle = 1.712$ Frame 34/89

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#### Random networks: examples for N=1000



#### **Poisson basics:**

Normalization: we must have

$$\sum_{k=0}^{\infty} P(k; \langle k \rangle) = 1$$

Checking:

$$\sum_{k=0}^{\infty} P(k; \langle k \rangle) = \sum_{k=0}^{\infty} \frac{\langle k \rangle^k}{k!} e^{-\langle k \rangle}$$
$$= e^{-\langle k \rangle} \sum_{k=0}^{\infty} \frac{\langle k \rangle^k}{k!}$$
$$= e^{-\langle k \rangle} e^{\langle k \rangle} = \mathbf{1} \checkmark$$

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#### Poisson basics:

Mean degree: we must have

$$\langle k \rangle = \sum_{k=0}^{\infty} k P(k; \langle k \rangle).$$

Checking:

$$\sum_{k=0}^{\infty} k P(k; \langle k \rangle) = \sum_{k=0}^{\infty} k \frac{\langle k \rangle^{k}}{k!} e^{-\langle k \rangle}$$
$$= e^{-\langle k \rangle} \sum_{k=1}^{\infty} \frac{\langle k \rangle^{k}}{(k-1)!}$$
$$= \langle k \rangle e^{-\langle k \rangle} \sum_{k=1}^{\infty} \frac{\langle k \rangle^{k-1}}{(k-1)!}$$
$$\langle k \rangle e^{-\langle k \rangle} \sum_{i=0}^{\infty} \frac{\langle k \rangle^{i}}{i!} = \langle k \rangle e^{-\langle k \rangle} e^{\langle k \rangle} = \langle k \rangle \checkmark$$

We'll get to a better way of doing this...

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# The edge-degree distribution:

- The degree distribution P<sub>k</sub> is fundamental for our description of many complex networks
- Again:  $P_k$  is the degree of randomly chosen node.
- A second very important distribution arises from choosing randomly on edges rather than on nodes.
- Define Q<sub>k</sub> to be the probability the node at a random end of a randomly chosen edge has degree k.
- Now choosing nodes based on their degree (i.e., size):
- Normalized form:

$$Q_k = \frac{kP_k}{\sum_{k'=0}^{\infty} k'P_{k'}} = \frac{kP_k}{\langle k \rangle}.$$

 $Q_k \propto k P_k$ 

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# Poisson basics:

- The variance of degree distributions for random networks turns out to be very important.
- Use calculation similar to one for finding (k) to find the second moment:

$$\langle k^2 \rangle = \langle k \rangle^2 + \langle k \rangle$$

Variance is then

$$\sigma^{2} = \langle \mathbf{k}^{2} \rangle - \langle \mathbf{k} \rangle^{2} = \langle \mathbf{k} \rangle^{2} + \langle \mathbf{k} \rangle - \langle \mathbf{k} \rangle^{2} = \langle \mathbf{k} \rangle.$$

- So standard deviation  $\sigma$  is equal to  $\sqrt{\langle k \rangle}$ .
- Note: This is a special property of Poisson distribution and can trip us up...

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# The edge-degree distribution:

- For random networks, Q<sub>k</sub> is also the probability that a friend (neighbor) of a random node has k friends.
- Useful variant on  $Q_k$ :

 $R_k$  = probability that a friend of a random node has *k* other friends.

$$R_k = rac{(k+1)P_{k+1}}{\sum_{k'=0}(k'+1)P_{k'+1}} = rac{(k+1)P_{k+1}}{\langle k 
angle}$$

- Equivalent to friend having degree k + 1.
- Natural question: what's the expected number of other friends that one friend has?

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# The edge-degree distribution:

Given R<sub>k</sub> is the probability that a friend has k other friends, then the average number of friends' other friends is

$$\langle k \rangle_R = \sum_{k=0}^{\infty} kR_k = \sum_{k=0}^{\infty} k \frac{(k+1)P_{k+1}}{\langle k \rangle}$$
$$= \frac{1}{\langle k \rangle} \sum_{k=1}^{\infty} k(k+1)P_{k+1}$$
$$= \frac{1}{\langle k \rangle} \sum_{k=1}^{\infty} \left( (k+1)^2 - (k+1) \right) P_{k+1}$$

(where we have sneakily matched up indices)

$$= \frac{1}{\langle k \rangle} \sum_{j=0}^{\infty} (j^2 - j) P_j \quad \text{(using j = k+1)}$$
$$= \frac{1}{\langle k \rangle} \left( \langle k^2 \rangle - \langle k \rangle \right)$$

#### Two reasons why this matters

Reason #1:

Average # friends of friends per node is

$$\langle k_2 \rangle = \langle k \rangle \times \langle k \rangle_R = \langle k \rangle \frac{1}{\langle k \rangle} \left( \langle k^2 \rangle - \langle k \rangle \right) = \langle k^2 \rangle - \langle k \rangle.$$

- Key: Average depends on the 1st and 2nd moments of P<sub>k</sub> and not just the 1st moment.
- Three peculiarities:
  - 1. We might guess  $\langle k_2 \rangle = \langle k \rangle (\langle k \rangle 1)$  but it's actually  $\langle k(k-1) \rangle$ .
  - 2. If  $P_k$  has a large second moment, then  $\langle k_2 \rangle$  will be big.
    - (e.g., in the case of a power-law distribution)
  - 3. Your friends are different to you...

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# The edge-degree distribution:

- Note: our result, ⟨k⟩<sub>R</sub> = 1/⟨k⟩ (⟨k²⟩ ⟨k⟩), is true for all random networks, independent of degree distribution.
- ▶ For standard random networks, recall

$$\langle k^2 \rangle = \langle k \rangle^2 + \langle k \rangle.$$

Therefore:

$$\left\langle k \right\rangle_{R} = \frac{1}{\left\langle k \right\rangle} \left( \left\langle k \right\rangle^{2} + \left\langle k \right\rangle - \left\langle k \right\rangle \right) = \left\langle k \right\rangle$$

- Again, neatness of results is a special property of the Poisson distribution.
- ► So friends on average have  $\langle k \rangle$  other friends, and  $\langle k \rangle$  + 1 total friends...

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### Two reasons why this matters

#### More on peculiarity #3:

- A node's average # of friends: (k)
- Friend's average # of friends:  $\frac{\langle k^2 \rangle}{\langle k \rangle}$
- Comparison:

$$\frac{\langle k^2 \rangle}{\langle k \rangle} = \langle k \rangle \frac{\langle k^2 \rangle}{\langle k \rangle^2} = \langle k \rangle \frac{\sigma^2 + \langle k \rangle^2}{\langle k \rangle^2} = \langle k \rangle \left( 1 + \frac{\sigma^2}{\langle k \rangle^2} \right) \ge$$

- So only if everyone has the same degree (variance= σ<sup>2</sup> = 0) can a node be the same as its friends.
- Intuition: for random networks, the more connected a node, the more likely it is to be chosen as a friend.

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#### Two reasons why this matters

#### (Big) Reason #2:

- \$\langle k \rangle\_R\$ is key to understanding how well random networks are connected together.
- e.g., we'd like to know what's the size of the largest component within a network.
- As N → ∞, does our network have a giant component?
- Defn: Component = connected subnetwork of nodes such that ∃ path between each pair of nodes in the subnetwork, and no node outside of the subnetwork is connected to it.
- Defn: Giant component = component that comprises a non-zero fraction of a network as N → ∞.
- Note: Component = Cluster

# Giant component

#### Standard random networks:

- Recall  $\langle k^2 \rangle = \langle k \rangle^2 + \langle k \rangle$ .
- Condition for giant component:

$$\langle k \rangle_{R} = \frac{\langle k^{2} \rangle - \langle k \rangle}{\langle k \rangle} = \frac{\langle k \rangle^{2} + \langle k \rangle - \langle k \rangle}{\langle k \rangle} = \langle k \rangle$$

- Therefore when (k) > 1, standard random networks have a giant component.
- When  $\langle k \rangle < 1$ , all components are finite.
- ► Fine example of a continuous phase transition (⊞).
- We say  $\langle k \rangle = 1$  marks the critical point of the system.



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### Structure of random networks

#### Giant component:

- A giant component exists if when we follow a random edge, we are likely to hit a node with at least 1 other outgoing edge.
- Equivalently, expect exponential growth in node number as we move out from a random node.
- All of this is the same as requiring  $\langle k \rangle_R > 1$ .
- Giant component condition (or percolation condition):

$$\langle k \rangle_R = rac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle} > 1$$

- Again, see that the second moment is an essential part of the story.
- Equivalent statement:  $\langle k^2 \rangle > 2 \langle k \rangle$

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# Giant component

Random networks with skewed  $P_k$ :

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• e.g, if 
$$P_k = ck^{-\gamma}$$
 with 2 <  $\gamma$  < 3 then

$$\langle k^2 
angle = c \sum_{k=0}^{\infty} k^2 k^{-\gamma}$$

$$\sim \int_{x=0}^{\infty} x^{2-\gamma} \mathrm{d}x$$
$$x^{3-\gamma}\Big|_{x=0}^{\infty} = \infty \quad (>\langle k \rangle).$$

- So giant component always exists for these kinds of networks.
- Cutoff scaling is k<sup>-3</sup>: if γ > 3 then we have to look harder at ⟨k⟩<sub>B</sub>.

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### Giant component

#### And how big is the largest component?

- Define  $S_1$  as the size of the largest component.
- Consider an infinite ER random network with average degree (k).
- Let's find  $S_1$  with a back-of-the-envelope argument.
- Define δ as the probability that a randomly chosen node does not belong to the largest component.
- Simple connection:  $\delta = 1 S_1$ .
- Dirty trick: If a randomly chosen node is not part of the largest component, then none of its neighbors are.
- So

$$\delta = \sum_{k=0}^{\infty} P_k \delta^k$$

Substitute in Poisson distribution...

# Giant component

- We can figure out some limits and details for  $S_1 = 1 e^{-\langle k \rangle S_1}$ .
- First, we can write  $\langle k \rangle$  in terms of  $S_1$ :

$$\langle k \rangle = rac{1}{S_1} \ln rac{1}{1-S_1}$$

- ▶ As  $\langle k \rangle \rightarrow 0$ ,  $S_1 \rightarrow 0$ .
- As  $\langle k \rangle \rightarrow \infty$ ,  $S_1 \rightarrow 1$ .
- Notice that at  $\langle k \rangle = 1$ , the critical point,  $S_1 = 0$ .
- Only solvable for S > 0 when  $\langle k \rangle > 1$ .
- ▶ Really a transcritical bifurcation<sup>[2]</sup>.

# Giant component

Carrying on:

$$\delta = \sum_{k=0}^{\infty} P_k \delta^k = \sum_{k=0}^{\infty} \frac{\langle k \rangle^k}{k!} e^{-\langle k \rangle} \delta^k$$
$$= e^{-\langle k \rangle} \sum_{k=0}^{\infty} \frac{(\langle k \rangle \delta)^k}{k!}$$
$$= e^{-\langle k \rangle} e^{\langle k \rangle \delta} = e^{-\langle k \rangle (1-\delta)}.$$

▶ Now substitute in  $\delta = 1 - S_1$  and rearrange to obtain:

$$S_1 = 1 - e^{-\langle k \rangle S_1}$$

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# Random Networks

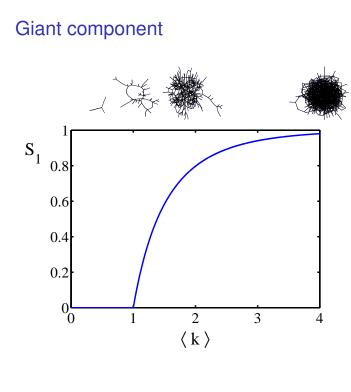
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### Giant component

#### Turns out we were lucky ...

- Our dirty trick only works for ER random networks.
- The problem: We assumed that neighbors have the same probability δ of belonging to the largest component.
- But we know our friends are different from us...
- Works for ER random networks because  $\langle k \rangle = \langle k \rangle_R$ .
- We need a separate probability δ' for the chance that a node at the end of a random edge is part of the largest component.
- We can do this but we need to enhance our toolkit with Generatingfunctionology...<sup>[3]</sup>



# Generating functions

- Idea: Given a sequence a<sub>0</sub>, a<sub>1</sub>, a<sub>2</sub>,..., associate each element with a distinct function or other mathematical object.
- Well-chosen functions allow us to manipulate sequences and retrieve sequence elements.

#### Definition:

• The generating function (g.f.) for a sequence  $\{a_n\}$  is

 $F(x)=\sum_{n=0}^{\infty}a_nx^n.$ 

- ► Roughly: transforms a vector in R<sup>∞</sup> into a function defined on R<sup>1</sup>.
- Related to Fourier, Laplace, Mellin, ...

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Rolling dice:

Simple example

$$F^{(\Box)}(x) = \sum_{k=1}^{6} p_k x^k = \frac{1}{6} (x + x^2 + x^3 + x^4 + x^5 + x^6)$$

We'll come back to this simple example as we derive various delicious properties of generating functions.



#### Example

Take a degree distribution with exponential decay:

 $P_k = ce^{-\lambda k}$ 

where  $c = 1 - e^{-\lambda}$ .

> The generating function for this distribution is

$$F(x) = \sum_{k=0}^{\infty} P_k x^k = \sum_{k=0}^{\infty} c e^{-\lambda k} x^k = \frac{c}{1 - x e^{-\lambda}}$$

- Notice that  $F(1) = c/(1 e^{-\lambda}) = 1$ .
- For probability distributions, we must always have
   F(1) = 1 since

$$F(1) = \sum_{k=0}^{\infty} P_k 1^k = \sum_{k=0}^{\infty} P_k = 1.$$

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## Properties of generating functions

► Average degree:

$$\langle k \rangle = \sum_{k=0}^{\infty} k P_k = \sum_{k=0}^{\infty} k P_k x^{k-1} \Big|_{x=1}$$
  
=  $\frac{d}{dx} F(x) \Big|_{x=1} = F'(1)$ 

- In general, many calculations become simple, if a little abstract.
- ► For our exponential example:

$$F'(x)=\frac{(1-e^{-\lambda})e^{-\lambda}}{(1-xe^{-\lambda})^2}.$$

So:

$$\langle k 
angle = F'(1) = rac{e^{-\lambda}}{(1-e^{-\lambda})}$$

# Edge-degree distribution

Recall our condition for a giant component:

$$\langle k \rangle_R = \frac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle} > 1.$$

- Let's r
   r
   express our condition in terms of generating functions.
- We first need the g.f. for  $R_k$ .
- ► We'll now use this notation:

 $F_P(x)$  is the g.f. for  $P_k$ .  $F_B(x)$  is the g.f. for  $R_k$ .

Condition in terms of g.f. is:

$$\langle k \rangle_R = F'_R(1) > 1$$

• Now find how  $F_R$  is related to  $F_P$ ...

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# Properties of generating functions

Useful pieces for probability distributions:

- Normalization:
- First moment:

$$\langle k \rangle = F'(1)$$

F(1) = 1

Higher moments:

$$\langle k^n \rangle = \left( x \frac{\mathrm{d}}{\mathrm{d}x} \right)^n F(x) \Big|_{x=1}$$

kth element of sequence (general):

 $P_k = \frac{1}{k!} \frac{\mathrm{d}^k}{\mathrm{d}x^k} F(x) \bigg|_{x=0}$ 

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#### Edge-degree distribution

We have

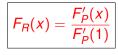
$$F_R(x) = \sum_{k=0}^{\infty} \mathbf{R}_k x^k = \sum_{k=0}^{\infty} \frac{(k+1)\mathbf{P}_{k+1}}{\langle k \rangle} x^k.$$

Shift index to j = k + 1 and pull out  $\frac{1}{\langle k \rangle}$ :

$$F_{R}(x) = \frac{1}{\langle k \rangle} \sum_{j=1}^{\infty} j P_{j} x^{j-1} = \frac{1}{\langle k \rangle} \sum_{j=1}^{\infty} P_{j} \frac{\mathrm{d}}{\mathrm{d}x} x^{j}$$

$$=\frac{1}{\langle k\rangle}\frac{\mathrm{d}}{\mathrm{d}x}\sum_{j=1}^{\infty}P_{j}x^{j}=\frac{1}{\langle k\rangle}\frac{\mathrm{d}}{\mathrm{d}x}\left(F_{P}(x)-P_{0}\right)=\frac{1}{\langle k\rangle}F_{P}'(x)$$

Finally, since  $\langle k \rangle = F'_P(1)$ ,



### Edge-degree distribution

- Recall giant component condition is  $\langle k \rangle_R = F'_R(1) > 1.$
- Since we have  $F_R(x) = F'_P(x)/F'_P(1)$ ,

$$F'_{R}(x) = rac{F''_{P}(x)}{F'_{P}(1).}$$

• Setting x = 1, our condition becomes



### Size distributions

G.f.'s for component size distributions:

$$F_{\pi}(x) = \sum_{n=0}^{\infty} \pi_n x^n$$
 and  $F_{\rho}(x) = \sum_{n=0}^{\infty} \rho_n x^n$ 

The largest component:

Subtle key:  $F_{\pi}(1)$  is the probability that a node belongs to a finite component.

• Therefore:  $S_1 = 1 - F_{\pi}(1)$ .

#### Our mission, which we accept:

Find the four generating functions

$$F_P, F_R, F_\pi$$
, and  $F_\rho$ .



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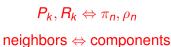
# Size distributions

To figure out the size of the largest component  $(S_1)$ , we need more resolution on component sizes.

#### **Definitions:**

- *π<sub>n</sub>* = probability that a random node belongs to a finite component of size *n* < ∞.</p>
- *ρ<sub>n</sub>* = probability a random link leads to a finite subcomponent of size *n* < ∞.</p>

#### Local-global connection:



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# Useful results we'll need for g.f.'s

#### Sneaky Result 1:

- Consider two random variables U and V whose values may be 0, 1, 2, ...
- Write probability distributions as U<sub>k</sub> and U<sub>k</sub> and g.f.'s as F<sub>U</sub> and F<sub>V</sub>.
- SR1: If a third random variable is defined as

$$W = \sum_{i=1}^{V} U^{(i)}$$
 with each  $U^{(i)} \stackrel{d}{=} U$ 

then

 $F_W(x) = F_V(F_U(x))$ 

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### Proof of SR1:

Write probability that variable W has value k as  $W_k$ .

$$W_{k} = \sum_{j=0}^{\infty} V_{j} \times \Pr(\text{sum of } j \text{ draws of variable } U = k)$$

$$= \sum_{j=0}^{\infty} V_{j} \sum_{\substack{\{i_{1}, i_{2}, \dots, i_{j}\} \mid \\ i_{1}+i_{2}+\dots+i_{j}=k}} U_{i_{1}} U_{i_{2}} \cdots U_{i_{j}}$$

$$\therefore F_{W}(x) = \sum_{k=0}^{\infty} W_{k} x^{k} = \sum_{k=0}^{\infty} \sum_{j=0}^{\infty} V_{j} \sum_{\substack{\{i_{1}, i_{2}, \dots, i_{j}\} \mid \\ i_{1}+i_{2}+\dots+i_{j}=k}} U_{i_{1}} U_{i_{2}} x^{i_{2}} \cdots U_{i_{j}} x^{i_{j}}}$$

$$= \sum_{j=0}^{\infty} V_{j} \sum_{k=0}^{\infty} \sum_{\substack{\{i_{1}, i_{2}, \dots, i_{j}\} \mid \\ i_{1}+i_{2}+\dots+i_{j}=k}} U_{i_{1}} x^{i_{1}} U_{i_{2}} x^{i_{2}} \cdots U_{i_{j}} x^{i_{j}}}$$

# Useful results we'll need for g.f.'s

Sneaky Result 2:

- Start with a random variable U with distribution  $U_k$  (k = 0, 1, 2, ...)
- SR2: If a second random variable is defined as

$$V = U + 1$$
 then  $F_V(x) = xF_U(x)$ 

• Reason:  $V_k = U_{k-1}$  for  $k \ge 1$  and  $V_0 = 0$ .

$$\therefore F_V(x) = \sum_{k=0}^{\infty} V_k x^k = \sum_{k=1}^{\infty} U_{k-1} x^k$$
$$= x \sum_{j=0}^{\infty} U_j x^j = x F_U(x) \cdot \checkmark$$

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# Proof of SR1:

With some concentration, observe:

Useful results we'll need for g.f.'s

Generalization of SR2:

▶ (1) If V = U + i then

▶ (2) If V = U - i then

$$F_{W}(x) = \sum_{j=0}^{\infty} V_{j} \sum_{k=0}^{\infty} \sum_{\substack{\{i_{1},i_{2},...,i_{j}\} \mid \\ i_{1}+i_{2}+...+i_{j}=k}} U_{i_{1}} x^{i_{1}} U_{i_{2}} x^{i_{2}} \cdots U_{i_{j}} x^{i_{j}}}{x^{k} \text{ piece of } \left(\sum_{i'=0}^{\infty} U_{i'} x^{i'}\right)^{j}} \left(\sum_{i'=0}^{\infty} U_{i'} x^{i'}\right)^{j} = (F_{U}(x))^{j}}{\left(\sum_{j=0}^{\infty} V_{j} (F_{U}(x))^{j}\right)} = F_{V} (F_{U}(x))^{j}$$

 $F_V(x) = x^i F_U(x).$ 

 $F_V(x) = x^{-i}F_U(x)$ 

 $= x^{-i} \sum_{k=0}^{\infty} U_k x^k$ 

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### Connecting generating functions

- Goal: figure out forms of the component generating functions, *F<sub>π</sub>* and *F<sub>ρ</sub>*.
- $\pi_n$  = probability that a random node belongs to a finite component of size *n*

 $= \sum_{k=0}^{\infty} P_k \times \Pr\left(\begin{array}{c} \text{sum of sizes of subcomponents} \\ \text{at end of } k \text{ random links} = n-1 \end{array}\right)$ 

Therefore:  $F_{\pi}(x) = \underbrace{x}_{SR2} \underbrace{F_{P}(F_{\rho}(x))}_{SR1}$ 

Extra factor of x accounts for random node itself.

# Connecting generating functions

We now have two functional equations connecting our generating functions:

 $F_{\pi}(x) = xF_P(F_{\rho}(x))$  and  $F_{\rho}(x) = xF_R(F_{\rho}(x))$ 

- Taking stock: We know  $F_P(x)$  and  $F_R(x) = F'_P(x)/F'_P(1)$ .
- We first untangle the second equation to find  $F_{\rho}$
- We can do this because it only involves  $F_{\rho}$  and  $F_{R}$ .
- The first equation then immediately gives us F<sub>π</sub> in terms of F<sub>ρ</sub> and F<sub>R</sub>.



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# Connecting generating functions

- $\rho_n$  = probability that a random link leads to a finite subcomponent of size *n*.
- Invoke one step of recursion: ρ<sub>n</sub> = probability that a random node arrived along a random edge is part of a finite subcomponent of size n.

$$= \sum_{k=0}^{\infty} R_k \times \Pr\left(\begin{array}{c} \text{sum of sizes of subcomponents} \\ \text{at end of } k \text{ random links} = n-1 \end{array}\right)$$

Therefore: 
$$F_{\rho}(x) = \underbrace{x}_{SR2} \underbrace{F_{R}(F_{\rho}(x))}_{SR1}$$

 Again, extra factor of x accounts for random node itself. Frame 79/89

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#### **Component sizes**

Remembering vaguely what we are doing:

Finding  $F_{\pi}$  to obtain the fractional size of the largest component  $S_1 = 1 - F_{\pi}(1)$ .

• Set x = 1 in our two equations:

 $F_{\pi}(1) = F_{P}(F_{\rho}(1))$  and  $F_{\rho}(1) = F_{R}(F_{\rho}(1))$ 

- Solve second equation numerically for  $F_{\rho}(1)$ .
- Plug  $F_{\rho}(1)$  into first equation to obtain  $F_{\pi}(1)$ .

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#### **Component sizes**

Example: Standard random graphs.

• We can show  $F_P(x) = e^{-\langle k \rangle (1-x)}$ 

. 
$$F_R(x) = F'_P(x)/F'_P(1) = e^{-\langle k \rangle (1-x)}/e^{-\langle k \rangle (1-x')}|_{x'=1}$$

 $= e^{-\langle k \rangle (1-x)} = F_P(x)$  ...aha!

- RHS's of our two equations are the same.
- So  $F_{\pi}(x) = F_{\rho}(x) = xF_R(F_{\rho}(x)) = xF_R(F_{\pi}(x))$
- Why our dirty (but wrong) trick worked earlier...

### Average component size

- Next: find average size of finite components  $\langle n \rangle$ .
- Using standard G.F. result:  $\langle n \rangle = F'_{\pi}(1)$ .
- Try to avoid finding  $F_{\pi}(x)$ ...
- ▶ Starting from  $F_{\pi}(x) = xF_{P}(F_{\rho}(x))$ , we differentiate:

$$F'_{\pi}(x) = F_{\mathcal{P}}\left(F_{\rho}(x)
ight) + xF'_{
ho}(x)F'_{\mathcal{P}}\left(F_{
ho}(x)
ight)$$

• While  $F_{\rho}(x) = xF_R(F_{\rho}(x))$  gives

 $F_{
ho}'(x)=F_R\left(F_{
ho}(x)
ight)+xF_{
ho}'(x)F_R'\left(F_{
ho}(x)
ight)$ 

- Now set x = 1 in both equations.
- We solve the second equation for F'<sub>ρ</sub>(1) (we must already have F<sub>ρ</sub>(1)).
- Plug  $F'_{\rho}(1)$  and  $F_{\rho}(1)$  into first equation to find  $F'_{\pi}(1)$ .

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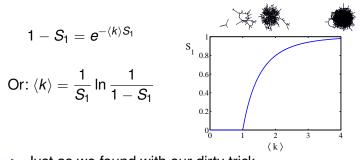
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#### **Component sizes**

- We are down to  $F_{\pi}(x) = xF_{R}(F_{\pi}(x))$  and  $F_{R}(x) = e^{-\langle k \rangle (1-x)}$ . •  $\therefore F_{\pi}(x) = xe^{-\langle k \rangle (1-F_{\pi}(x))}$
- We're first after S<sub>1</sub> = 1 − F<sub>π</sub>(1) so set x = 1 and replace F<sub>π</sub>(1) by 1 − S<sub>1</sub>:



Just as we found with our dirty trick ...
Again, we (usually) have to resort to numerics ...

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End result:  $\langle n \rangle = F'_{\pi}(1) = \frac{(1-S_1)}{1-\langle k \rangle(1-S_1)}$ 

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# Average component size

Example: Standard random graphs.

- Use fact that  $F_P = F_R$  and  $F_{\pi} = F_{\rho}$ .
- Two differentiated equations reduce to only one:

 $F'_{\pi}(x) = F_P\left(F_{\pi}(x)\right) + xF'_{\pi}(x)F'_P\left(F_{\pi}(x)\right)$ 

Rearrange: 
$$F'_{\pi}(x) = \frac{F_P(F_{\pi}(x))}{1 - xF'_P(F_{\pi}(x))}$$

Simplify denominator using 
$$F'_P(x) = \langle k \rangle F_P(x)$$

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#### Average component size

Our result for standard random networks:

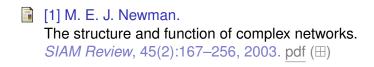
$$|n\rangle = F'_{\pi}(1) = rac{(1-S_1)}{1-\langle k 
angle(1-S_1)}$$

- Recall that (k) = 1 is the critical value of average degree for standard random networks.
- Look at what happens when we increase (k) to 1 from below.
- We have  $S_1 = 0$  for all  $\langle k \rangle < 1$  so

$$\langle n \rangle = \frac{1}{1 - \langle k \rangle}$$

- This blows up as  $\langle k \rangle \rightarrow 1$ .
- Reason: we have a power law distribution of component sizes at (k) = 1.
- Typical critical point behavior....

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#### Average component size

• Limits of  $\langle k \rangle = 0$  and  $\infty$  make sense for

$$\langle n \rangle = F'_{\pi}(1) = rac{(1-S_1)}{1-\langle k \rangle(1-S_1)}$$

- As  $\langle k \rangle \rightarrow 0$ ,  $S_1 = 0$ , and  $\langle n \rangle \rightarrow 1$ .
- All nodes are isolated.
- As  $\langle k \rangle \rightarrow \infty$ ,  $S_1 \rightarrow 1$  and  $\langle n \rangle \rightarrow 0$ .
- No nodes are outside of the giant component.

#### Extra on largest component size:

- For  $\langle k \rangle = 1$ ,  $S_1 \sim N^{2/3}$ .
- For  $\langle k \rangle < 1$ ,  $S_1 \sim \log N$ .

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