Random Networks Complex Networks, Course 295A, Spring, 2008

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Random networks

Pure, abstract random networks:

- Consider set of all networks with N labelled nodes and m edges.
- Standard random network = randomly chosen network from this set.
- To be clear: each network is equally probable.
- Sometimes equiprobability is a good assumption, but it is always an assumption.
- Known as Erdös-Rényi random networks or ER graphs.

Outline

Random Networks

Basics

Structure

Frame 1/83

B 990

Random Networks

Basics

Definitions

Basics

Definitions How to build Some visual examples

Structure

Clustering Degree distributions Configuration model Largest component

Generating Functions Definitions

Properties

References

Basics Definitions How to build Some visual examples Structure Clustering

Random Network

Configuration mod Largest componer

> Definitions Properties

References

Random Network

Basics

Definitions

Structure

Generating

Random networks

Some features:

networks.

Number of possible edges:

$$0 \le m \le \binom{N}{2} = \frac{N(N-1)}{2}$$

- Given *m* edges, there are $\binom{\binom{N}{2}}{m}$ different possible
- Crazy factorial explosion for $1 \ll m \ll \binom{N}{2}$.
- Limit of m = 0: empty graph.
- Limit of $m = \binom{N}{2}$: complete or fully-connected graph.
- Real world: links are usually costly so real networks are almost always sparse.

Random networks

How to build standard random networks:

- ▶ Given N and m.
- Two probablistic methods (we'll see a third later on)
- 1. Connect each of the $\binom{N}{2}$ pairs with appropriate probability *p*.
 - Useful for theoretical work.
- 2. Take *N* nodes and add exactly *m* links by selecting edges without replacement.
 - ► Algorithm: Randomly choose a pair of nodes *i* and *j*, *i* ≠ *j*, and connect if unconnected; repeat until all *m* edges are allocated.
 - Best for adding small numbers of links (most cases).
 - ▶ 1 and 2 are effectively equivalent for large *N*.



Random networks

A few more things:

For method 1, # links is probablistic:

$$\langle m \rangle = p \binom{N}{2} = p \frac{1}{2} N(N-1)$$

► So the expected or average degree is

$$\langle k \rangle = \frac{2 \langle m \rangle}{N}$$
$$= \frac{2}{N} \rho \frac{1}{2} N(N-1) = \frac{2}{N} \rho \frac{1}{2} N(N-1) = \rho(N-1)$$

- Which is what it should be...
- If we keep $\langle k \rangle$ constant then $p \propto 1/N \to 0$ as $N \to \infty$.

How to build Some Visual example Structure Clustering Degree distributions Configuration model Largest component Generating Functions Properties

Random Networks

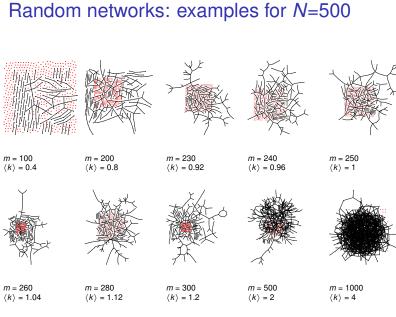
Basics

Random Networks

Some visual example:

Generating

Basics

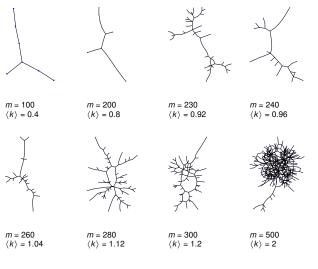


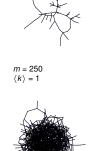


Frame 7/83

B 990

Random networks: largest components



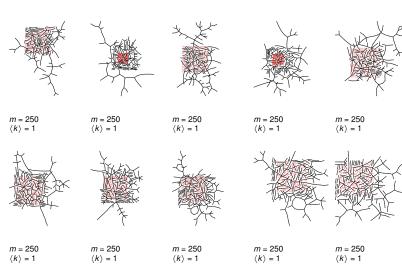


m = 1000

 $\langle k \rangle = 4$

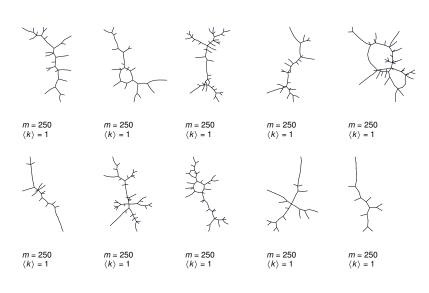
Frame 22/83

Random networks: examples for N=500





Random networks: largest components



Frame 24/83

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Random Networks

Some visual example:

Generating

Basics

Random networks

Clustering:

- For method 1, what is the clustering coefficient for a finite network?
- Consider triangle/triple clustering coefficient (Newman^[1]):

$$C_2 = \frac{3 \times \# triangles}{\# triples}$$

- Recall: C₂ = probability that two nodes are connected given they have a friend in common.
- > For standard random networks, we have simply that

$C_2 = p.$

Random Networks
Basics
Definitions
Some visual examples
Structure
Clustering
Generating
Functions
References

Frame 26/83

P

Random networks

Clustering:

- So for large random networks (N → ∞), clustering drops to zero.
- Key structural feature of random networks is that they locally look like branching networks (no loops).

Random Networks

Basics Definitions How to build Some visual examples Structure Clustering Degree distributions Configuration model

Generating Functions Definitions Properties

References

Frame 27/83

Random networks

Degree distribution:

- Recall p_k = probability that a randomly selected node has degree k.
- Consider method 1 for constructing random networks: each possible link is realized with probability p.
- Now consider one node: there are 'N choose k' ways the node can be connected to k of the other N – 1 nodes.
- Each connection occurs with probability p, each non-connection with probability (1 – p).
- Therefore have a binomial distribution:

$$P(k; p, N) = \binom{N-1}{k} p^k (1-p)^{N-1-k}$$

Limiting	form	of	P (<i>k</i> :	D.	N):
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Substitute
$$p = \frac{\langle k \rangle}{N-1}$$
 into $P(k; p, N)$ and hold k fixed:

$$P(k; p, N) = \binom{N-1}{k} \left(\frac{\langle k \rangle}{N-1} \right)^k \left(1 - \frac{\langle k \rangle}{N-1} \right)^{N-1-k}$$

$$= \frac{(N-1)!}{k! (N-1-k)!} \frac{\langle k \rangle^k}{(N-1)^k} \left(1 - \frac{\langle k \rangle}{N-1} \right)^{N-1-k}$$

$$= \frac{(N-1)(N-2)\cdots(N-k)}{k!} \frac{\langle k \rangle^k}{(N-1)^k} \left(1 - \frac{\langle k \rangle}{N-1} \right)^{N-1-k}$$

Random Networks

Basics

Degree distributions

Frame 29/83

B 990

Random networks

Limiting form of P(k; p, N):

- Our degree distribution: $P(k; p, N) = \binom{N-1}{k} p^k (1-p)^{N-1-k}.$
- What happens as $N \to \infty$?
- We must end up with the normal distribution right?
- If *p* is fixed, then we would end up with a Gaussian with average degree ⟨*k*⟩ ≃ *pN* → ∞.
- But we want to keep $\langle k \rangle$ fixed...
- ► So examine limit of P(k; p, N) when $p \to 0$ and $N \to \infty$ with $\langle k \rangle = p(N 1) = \text{constant}$.



Random Network

Frame 30/83

Random Network

Limiting form of P(k; p, N):

► We are now here:

$$P(k; p, N) \simeq rac{\langle k
angle^k}{k!} \left(1 - rac{\langle k
angle}{N-1}
ight)^{N-1-k}$$

Now use the excellent result:

$$\lim_{n\to\infty}\left(1+\frac{x}{n}\right)^n=e^x$$

(Use l'Hôpital's rule to prove.)

• Identifying n = N - 1 and $x = -\langle k \rangle$:

$$\mathcal{P}(k;\langle k
angle)\simeqrac{\langle k
angle^k}{k!}m{e}^{-\langle k
angle}\left(1-rac{\langle k
angle}{N-1}
ight)^{-k}
ightarrowrac{\langle k
angle^k}{k!}m{e}^{-\langle k
angle}$$

▶ This is a Poisson distribution (\boxplus) with mean $\langle k \rangle$.

Basics Definitions How to build Some visual exampl Structure Clustering Degree distributions Configuration model Largest component Generating

> Definitions Properties

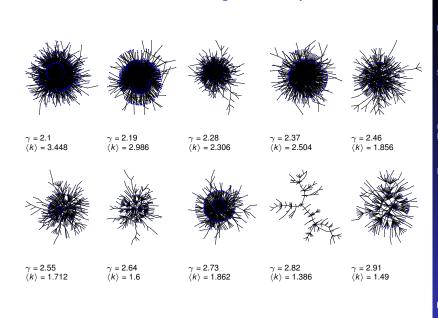
References

General random networks

- So... standard random networks have a Poisson degree distribution
- Generalize to arbitrary degree distribution P_k .
- Also known as the configuration model^[1].
- Can generalize construction method from ER random networks.
- Assign each node a weight w from some distribution P_w and form links with probability

 $P(\text{link between } i \text{ and } j) \propto w_i w_j.$

- But we'll be more interested in
 - 1. Randomly wiring up (and rewiring) already existing nodes with fixed degrees.
 - 2. Examining mechanisms that lead to networks with certain degree distributions.



Random networks: largest components



Poisson basics:

Checking:

Normalization: we must have

Random Networks

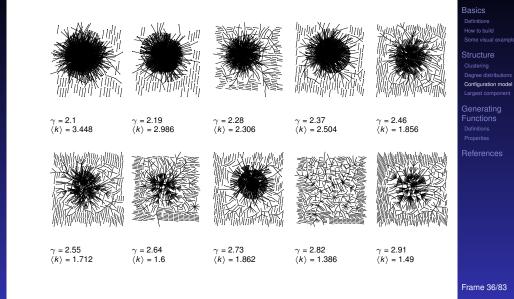
Basics

Configuration mode

Frame 34/83

B 9900

Random networks: examples for N=1000



 $\sum_{k=0}^{\infty} P(k; \langle k \rangle) = 1$

 $\sum_{k=0}^{\infty} P(k; \langle k \rangle) = \sum_{k=0}^{\infty} \frac{\langle k \rangle^k}{k!} e^{-\langle k \rangle}$

 $= e^{-\langle k
angle} \sum_{k=0}^{\infty} rac{\langle k
angle^k}{k!}$

 $= e^{-\langle k \rangle} e^{\langle k \rangle} = 1 \checkmark$

Basics Definitions How to build Some visual examples Structure Clustering Degree distributions Configuration model

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Random Network

Random Networks

Generating Functions Definitions Properties

Reference

Frame 38/83

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Poisson basics:

Mean degree: we must have

$$\langle k
angle = \sum_{k=0}^{\infty} k P(k; \langle k
angle)$$

Checking:

$$\sum_{k=0}^{\infty} kP(k; \langle k \rangle) = \sum_{k=0}^{\infty} k \frac{\langle k \rangle^{k}}{k!} e^{-\langle k \rangle}$$
$$= e^{-\langle k \rangle} \sum_{k=1}^{\infty} \frac{\langle k \rangle^{k}}{(k-1)!}$$
$$= \langle k \rangle e^{-\langle k \rangle} \sum_{k=1}^{\infty} \frac{\langle k \rangle^{k-1}}{(k-1)!}$$
$$\langle k \rangle e^{-\langle k \rangle} \sum_{i=0}^{\infty} \frac{\langle k \rangle^{i}}{i!} = \langle k \rangle e^{-\langle k \rangle} e^{\langle k \rangle} = \langle k \rangle \checkmark$$

We'll get to a better way of doing this...

=

The edge-degree distribution:

- The degree distribution P_k is fundamental for our description of many complex networks
- ► Again: *P_k* is the degree of randomly chosen node.
- A second very important distribution arises from choosing randomly on edges rather than on nodes.
- Define Q_k to be the probability the node at a random end of a randomly chosen edge has degree k.
- Now choosing nodes based on their degree (i.e., size):

 $Q_k \propto k P_k$

Normalized form:

$$Q_{k} = \frac{kP_{k}}{\sum_{k'=0}^{\infty} k'P_{k'}} = \frac{kP_{k}}{\langle k \rangle}.$$

Random Networks

Random Networks

Basics

Configuration mode

Generating

Frame 41/83

Poisson basics:

- The variance of degree distributions for random networks turns out to be very important.
- Use calculation similar to one for finding (k) to find the second moment:

$$\langle k^2 \rangle = \langle k \rangle^2 + \langle k \rangle$$

Variance is then

$$\sigma^{2} = \langle \mathbf{k}^{2} \rangle - \langle \mathbf{k} \rangle^{2} = \langle \mathbf{k} \rangle^{2} + \langle \mathbf{k} \rangle - \langle \mathbf{k} \rangle^{2} = \langle \mathbf{k} \rangle.$$

- So standard deviation σ is equal to $\sqrt{\langle k \rangle}$.
- Note: This is a special property of Poisson distribution and can trip us up...

Random Network

Basics

Configuration mode

Generating

Random Networks

Basics

Structure

Configuration mode

Generating

The edge-degree distribution:

- For random networks, Q_k is also the probability that a friend (neighbor) of a random node has k friends.
- Useful variant on Q_k :

 R_k = probability that a friend of a random node has *k* other friends.

$$\mathbf{R}_{k} = \frac{(k+1)P_{k+1}}{\sum_{k'=0}(k'+1)P_{k'+1}} = \frac{(k+1)P_{k+1}}{\langle k \rangle}$$

- Equivalent to friend having degree k + 1.
- Natural question: what's the expected number of other friends that one friend has?

Frame 42/83

The edge-degree distribution:

Given R_k is the probability that a friend has k other friends, then the average number of friends' other friends is

$$\begin{split} \langle k \rangle_{R} &= \sum_{k=0}^{\infty} k R_{k} = \sum_{k=0}^{\infty} k \frac{(k+1) P_{k+1}}{\langle k \rangle} \\ &= \frac{1}{\langle k \rangle} \sum_{k=1}^{\infty} k(k+1) P_{k+1} \\ &= \frac{1}{\langle k \rangle} \sum_{k=1}^{\infty} \left((k+1)^{2} - (k+1) \right) P_{k+1} \end{split}$$

(where we have sneakily matched up indices)

$$= \frac{1}{\langle k \rangle} \sum_{j=0}^{\infty} (j^2 - j) P_j \quad \text{(using j = k+1)}$$
$$= \frac{1}{\langle k \rangle} \left(\langle k^2 \rangle - \langle k \rangle \right)$$

Two reasons why this matters

Reason #1:

Average # friends of friends per node is

$$\langle k_2 \rangle = \langle k \rangle \times \langle k \rangle_R = \langle k \rangle \frac{1}{\langle k \rangle} \left(\langle k^2 \rangle - \langle k \rangle \right) = \langle k^2 \rangle - \langle k \rangle.$$

- Key: Average depends on the 1st and 2nd moments of P_k and not just the 1st moment.
- Three peculiarities:
 - 1. We might guess $\langle k_2 \rangle = \langle k \rangle (\langle k \rangle 1)$ but it's actually $\langle k(k-1) \rangle$.
 - 2. If P_k has a large second moment,
 - then $\langle k_2 \rangle$ will be big.
 - (e.g., in the case of a power-law distribution)
 - 3. Your friends are different to you...



- Note: our result, ⟨k⟩_R = 1/⟨k⟩ (⟨k²⟩ ⟨k⟩), is true for all random networks, independent of degree distribution.
- For standard random networks, recall

$$\langle k^2 \rangle = \langle k \rangle^2 + \langle k \rangle$$

► Therefore:

Random Networks

Basics

Structure

Configuration mode

Generating

Frame 43/83

B 990

Random Networks

Basics

Structure

Configuration mode

Generating Functions

Frame 4<u>5/83</u>

P

$$\langle k \rangle_{R} = \frac{1}{\langle k \rangle} \left(\langle k \rangle^{2} + \langle k \rangle - \langle k \rangle \right) = \langle k \rangle$$

- Again, neatness of results is a special property of the Poisson distribution.
- So friends on average have $\langle k \rangle$ other friends, and $\langle k \rangle + 1$ total friends...

Frame 44/83 日 のへへ

Random Network

Basics

Two reasons why this matters

More on peculiarity #3:

- A node's average # of friends: (k)
- Friend's average # of friends: $\frac{\langle k^2 \rangle}{\langle k \rangle}$
- Comparison:

$$\frac{\langle k^2 \rangle}{\langle k \rangle} = \langle k \rangle \frac{\langle k^2 \rangle}{\langle k \rangle^2} = \langle k \rangle \frac{\sigma^2 + \langle k \rangle^2}{\langle k \rangle^2} = \langle k \rangle \left(1 + \frac{\sigma^2}{\langle k \rangle^2} \right) \ge \langle k \rangle$$

- So only if everyone has the same degree (variance= σ² = 0) can a node be the same as its friends.
- Intuition: for random networks, the more connected a node, the more likely it is to be chosen as a friend.

Some visual examples Structure Clustering Degree distributions Configuration model Largest component Generating Functions Definitions Properties

Basics

Structure

Configuration mode

Generating

Two reasons why this matters

(Big) Reason #2:

- \$\langle k \rangle_R\$ is key to understanding how well random networks are connected together.
- e.g., we'd like to know what's the size of the largest component within a network.
- As N → ∞, does our network have a giant component?
- Defn: Component = connected subnetwork of nodes such that ∃ path between each pair of nodes in the subnetwork, and no node out side of the subnetwork is connected to it.
- Defn: Giant component = component that comprises a non-zero fraction of a network as N → ∞.
- Note: Component = Cluster

Giant component

Standard random networks:

- Recall $\langle k^2 \rangle = \langle k \rangle^2 + \langle k \rangle$.
- Condition for giant component:

$$\langle k \rangle_R = \frac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle} = \frac{\langle k \rangle^2 + \langle k \rangle - \langle k \rangle}{\langle k \rangle} = \langle k \rangle$$

- Therefore when (k) > 1, standard random networks have a giant component.
- When $\langle k \rangle < 1$, all components are finite.
- Fine example of a continuous phase transition (\boxplus) .
- We say $\langle k \rangle = 1$ marks the critical point of the system.

Structure of random networks

Giant component:

- A giant component exists if when we follow a random edge, we are likely to hit a node with at least 1 other outgoing edge.
- Equivalently, expect exponential growth in node number as we move out from a random node.
- All of this is the same as requiring $\langle k \rangle_R > 1$.
- Giant component condition (or percolation condition):

$$\langle k
angle_R = rac{\langle k^2
angle - \langle k
angle}{\langle k
angle} > 1$$

- Again, see that the second moment is an essential part of the story.
- Equivalent statement: $\langle k^2 \rangle > 2 \langle k \rangle$

Random Network

Giant component

Random networks with skewed P_k :

• e.g, if $P_k = ck^{-\gamma}$ with 2 < γ < 3 then

$$\langle k^2
angle = c \sum_{k=0}^{\infty} k^2 k^{-\gamma}$$

 $\sim \int_{x=0}^{\infty} x^{2-\gamma} dx$
 $\propto x^{3-\gamma} \Big|_{x=0}^{\infty} = \infty \quad (> \langle k \rangle).$

- So giant component always exists for these kinds of networks.
- Cutoff scaling is k⁻³: if γ > 3 then we have to look harder at ⟨k⟩_R.

Basics Definitions How to build Some visual exampl Structure Clustering Degree distributions Configuration model Largest component Generating

> Functions Definitions Properties

References

Basics Definitions How to build Some visual examples Structure Clustering Degree distributions Configuration model Largest component Generating Functions Primitions Properties References

Frame 50/83

Frame 47/83

B 990

Random Networks



Basics

Structure

Largest componen

Generating

unctions

Giant component

And how big is the largest component?

- Define S_1 as the size of the largest component.
- Consider an infinite ER random network with average degree (k).
- Let's find S_1 with a back-of-the-envelope argument.
- Define δ as the probability that a randomly chosen node does not belong to the largest component.
- Simple connection: $\delta = 1 S_1$.
- Dirty trick: If a randomly chosen node is not part of the largest component, then none of its neighbors are.
- So

$$\delta = \sum_{k=0}^{\infty} P_k \delta^k$$

Substitute in Poisson distribution...

Giant component

We can figure out some limits and details for

$$S_1 = 1 - e^{-\langle k \rangle S_1}.$$

- ▶ As $\langle k \rangle \rightarrow 0$, $S_1 \rightarrow 0$.
- ▶ As $\langle k \rangle \rightarrow \infty$, $S_1 \rightarrow 1$.
- Notice that at $\langle k \rangle = 1$, the critical point, $S_1 = 0$.
- Only solvable for S > 0 when $\langle k \rangle > 1$.
- Really a transcritical bifurcation^[2].



Random Networks

Basics

Structure

Largest component

Frame 52/83

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Random Networks

Basics

Structure

Largest component

Frame 54/83

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Carrying on:

$$\delta = \sum_{k=0}^{\infty} P_k \delta^k = \sum_{k=0}^{\infty} \frac{\langle k \rangle^k}{k!} e^{-\langle k \rangle} \delta^k$$
$$= e^{-\langle k \rangle} \sum_{k=0}^{\infty} \frac{(\langle k \rangle \delta)^k}{k!}$$
$$= e^{-\langle k \rangle} e^{\langle k \rangle \delta} = e^{-\langle k \rangle (1-\delta)}.$$

Now substitute in δ = 1 − S₁ and rearrange to obtain a transcendental equation for S₁:

$$S_1 = 1 - e^{-\langle k \rangle S_1}.$$

Frame 53/83

Random Networks

Basics

Structure

Largest component

Generating

Giant component

Random Network

Basics Definitions How to build Some visual example Structure

Degree distributions Configuration model Largest component

Generating Functions Definitions Properties

References

Frame 55/83

Giant component

Turns out we were lucky...

- Our dirty trick only works for ER random networks.
- The problem: We assumed that neighbors have the same probability δ of belonging to the largest component.
- But we know our friends are different from us...
- Works for ER random networks because $\langle k \rangle = \langle k \rangle_R$.
- We need a separate probability δ' for the chance that a node at the end of a random edge is part of the largest component.
- We can do this but we need to enhance our toolkit with Generatingfunctionology...^[3]

Random Networks
Basics
Definitions
How to build
Some visual examples
Structure
Clustering
Degree distributions
Configuration model
Largest component
Generating
Punctions
Properties
References

Frame 56/83

B 990

Random Networks

Generating functions

- Idea: Given a sequence a₀, a₁, a₂,..., associate each element with a distinct function or other mathematical object.
- Well-chosen functions allow us to manipulate sequences and retrieve sequence elements.

Definition:

• The generating function (g.f.) for a sequence $\{a_n\}$ is

$$F(x)=\sum_{n=0}^{\infty}a_nx^n.$$

- ► Roughly: transforms a vector in R[∞] into a function defined on R¹.
- ▶ Related to Fourier, Laplace, Mellin, ...

Random Network

Random Networks

Basics

Structure

Generating

Definitions

Example

Take a degree distribution with exponential decay:

 $P_k = ce^{-\lambda k}$

where $c = 1 - e^{-\lambda}$.

• The generating function for this distribution is

$$F(x) = \sum_{k=0}^{\infty} P_k x^k = \sum_{k=0}^{\infty} c e^{-\lambda k} x^k = \frac{c}{1 - x e^{-\lambda k}}$$

- Notice that $F(1) = c/(1 e^{-\lambda}) = 1$.
- For probability distributions, we must always have F(1) = 1 since

$$F(1) = \sum_{k=0}^{\infty} P_k 1^k = \sum_{k=0}^{\infty} P_k = 1$$

Basics Definitions How to build Some visual examples Structure Clustering Degree distributions Configuration model Largest component Generating Functions Properties References

Properties of generating functions

► Average degree:

$$\langle k \rangle = \sum_{k=0}^{\infty} k P_k = \sum_{k=0}^{\infty} k P_k x^{k-1} \bigg|_{x=1}$$
$$= \left. \frac{\mathrm{d}}{\mathrm{d}x} F(x) \right|_{x=1} = F'(1)$$

- In general, many calculations become simple, if a little astract.
- ► For our exponential example:

$$F'(x)=rac{(1-e^{-\lambda})e^{-\lambda}}{(1-xe^{-\lambda})^2}.$$

So:

$$\langle k \rangle = F'(1) = rac{e^{-\lambda}}{(1-e^{-\lambda})}.$$

Basics Definitions How to build Some visual examples Structure Clustering Degree distributions Configuration model Largest component Generating

Properties

Frame 61/83

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Properties of generating functions

Useful pieces for probability distributions:

Normalization:

F(1) = 1

First moment:

$$\langle k \rangle = F'(1)$$

Higher moments:

$$\langle k^n \rangle = \left(x \frac{\mathrm{d}}{\mathrm{d}x} \right)^n F(x) \Big|_{x=1}$$

kth element of sequence (general):

$$P_k = \frac{1}{k!} \frac{\mathrm{d}^k}{\mathrm{d}x^k} F(x) \bigg|_{x=0}$$

Edge-degree distribution

We have

$$F_{R}(x) = \sum_{k=0}^{\infty} \frac{R_{k}x^{k}}{k} = \sum_{k=0}^{\infty} \frac{(k+1)P_{k+1}}{\langle k \rangle} x^{k}$$

Shift index to j = k + 1 and pull out $\frac{1}{\langle k \rangle}$:

$$F_R(x) = \frac{1}{\langle k \rangle} \sum_{j=1}^{\infty} j P_j x^{j-1} = \frac{1}{\langle k \rangle} \sum_{j=0}^{\infty} P_j j x^{j-1}$$

$$= \frac{1}{\langle k \rangle} \sum_{j=0}^{\infty} \frac{\mathrm{d}}{\mathrm{d}x} P_j x^j = \frac{1}{\langle k \rangle} \frac{\mathrm{d}}{\mathrm{d}x} \sum_{j=0}^{\infty} P_j x^j = \frac{1}{\langle k \rangle} F_P'(x)$$

Finally, since $\langle k \rangle = F'_P(1)$,

 $F_R(x) = \frac{F'_P(x)}{F'_P(1)}$

Random Networks

Random Networks

Basics

Structure

Properties

Frame 64/83

P

Edge-degree distribution

Recall our condition for a giant component:

$$\langle k \rangle_{R} = \frac{\langle k^{2} \rangle - \langle k \rangle}{\langle k \rangle} > 1.$$

- We first need the g.f. for R_k .
- We'll now use this notation: $F_P(x)$ is the g.f. for P_k . $F_R(x)$ is the g.f. for R_k .
- Condition in terms of g.f. is:

$$\langle k \rangle_R = F'_R(1) > 1$$

• Now find how F_R is related to F_P ...

Frame 63/83 日 のへへ

Random Networks

Basics

Structure

Properties

Edge-degree distribution

- Recall giant component condition is $\langle k \rangle_R = F'_R(1) > 1.$
- Since we have $F_R(x) = F'_P(x)/F'_P(1)$,

$$F_R'(x) = \frac{F_P''(x)}{F_P'(1)}$$

• Setting x = 1, our condition becomes



Basics Definitions

Random Networks

Some visual example
Structure
Clustering

Configuration mode Largest component Generating

> -UNCTIONS Definitions Properties

References

Frame 65/83

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Size distributions

To figure out the size of the largest component (S_1) , we need more resolution on component sizes.

Definitions:

- $\blacktriangleright \pi_n$ = probability that a random node belongs to a finite component of size $n < \infty$.
- ρ_n = probability a random link leads to a finite subcomponent of size $n < \infty$.

Local-global connection:

 $P_k, R_k \Leftrightarrow \pi_n, \rho_n$ neighbors \Leftrightarrow components

Useful results we'll need for g.f.'s

Sneaky Result 1:

- Consider two random variables U and V whose values may be 0, 1, 2, ...
- Write probability distributions as U_k and V_k and g.f.'s as F_U and F_V .
- SR1: If a third random variable is defined as

$$W = \sum_{i=1}^{V} U^{(i)}$$
 with each $U^{(i)} \stackrel{d}{=} U$

then

$$F_W(x) = F_V(F_U(x))$$

Random Network
Basics Definitions How to build Some visual examples Structure Clustering Degree distributions Configuration model
Largest component Generating Functions Definitions Properties References
Frame 68/83
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Random Networks

Basics

Properties

Frame 66/83

B 990

Size distributions G.f.'s for component size distributions:

 $F_{\pi}(x) = \sum_{k=0}^{\infty} \pi_n x^n$ and $F_{\rho}(x) = \sum_{k=0}^{\infty} \rho_n x^n$

The largest component:

- Subtle key: $F_{\pi}(1)$ is the probability that a node belongs to a finite component.
- Therefore: $S_1 = 1 F_{\pi}(1)$.

Our mission, which we accept:

Find the four generating functions

 F_P, F_B, F_{π} , and F_o .

Proof of SN1:

Write probability that variable W has value k as W_k .

 $W_k = \sum_{i=0}^{\infty} V_j \times \Pr(\text{sum of } j \text{ draws of variable } U = k)$

$$= \sum_{j=0}^{\infty} V_j \sum_{\substack{\{i_1, i_2, \dots, i_k\} \mid \\ i_1+i_2+\dots+i_k=j}} U_{i_1} U_{i_2} \cdots U_{i_j}$$

$$\therefore F_{W}(x) = \sum_{k=0}^{\infty} W_{k} x^{k} = \sum_{k=0}^{\infty} \sum_{j=0}^{\infty} V_{j} \sum_{\substack{\{i_{1}, i_{2}, \dots, i_{k}\} \mid \\ i_{1}+i_{2}+\dots+i_{k}=j}}^{\sum} U_{i_{1}} U_{i_{2}} x^{i_{2}} \cdots U_{i_{j}} x^{i_{j}}}$$
$$= \sum_{j=0}^{\infty} V_{j} \sum_{k=0}^{\infty} \sum_{\substack{\{i_{1}, i_{2}, \dots, i_{k}\} \mid \\ \{i_{1}, i_{2}, \dots, i_{k}\} \mid }}^{\sum} U_{i_{1}} x^{i_{1}} U_{i_{2}} x^{i_{2}} \cdots U_{i_{j}} x^{i_{j}}}$$

$$\sum_{i_1,i_2,\dots,i_k\}|\atop{i_1+i_2+\dots+i_k=j}} O_{I_1} \times O_{I_2} \times O_{I_2}$$

Basics Generating Properties

Random Network

Frame 67/83 B 990

Basics

Random Network

unctions Properties

Frame 69/83 ð

Proof of SN1:

With some concentration, observe:

$$F_{W}(x) = \sum_{j=0}^{\infty} V_{j} \sum_{k=0}^{\infty} \sum_{\substack{\{i_{1}, i_{2}, \dots, i_{k}\} \mid \\ i_{1}+i_{2}+\dots+i_{k}=j}} U_{i_{1}} x^{i_{1}} U_{i_{2}} x^{i_{2}} \dots U_{i_{j}} x^{i_{j}}}{x^{k} \text{ piece of } \left(\sum_{i'=0}^{\infty} U_{i'} x^{i'}\right)^{j}} \left(\sum_{i'=0}^{\infty} U_{i'} x^{i'}\right)^{j}} = (F_{U}(x))^{j}}$$
$$= \sum_{j=0}^{\infty} V_{j} (F_{U}(x))^{j}$$
$$= F_{V} (F_{U}(x)) \checkmark$$

Useful results we'll need for g.f.'s

Generalization of SN2:

• (1) If V = U + i then

$$F_V(x) = x^i F_U(x).$$

▶ (2) If V = U - i then

$$F_V(x) = x^{-i} \left(F_U(x) - U_0 - U_1 x - \ldots - U_{i-1} x^{i-1} \right)$$

$$= x^{-i} \sum_{k=i}^{\infty} U_k x^k$$

Random Networks

Random Networks

Basics

Structure

Properties

Frame 72/83

B 990

Useful results we'll need for g.f.'s

Sneaky Result 2:

- Start with a random variable U with distribution U_k (k = 0, 1, 2, ...)
- SNR2: If a second random variable is defined as

V = U + 1 then $F_V(x) = xF_U(x)$

• Reason:
$$V_k = U_{k-1}$$
 for $k \ge 1$ and $V_0 = 0$

$$\therefore F_V(x) = \sum_{k=0}^{\infty} V_k x^k = \sum_{k=1}^{\infty} U_{k-1} x^k$$
$$= x \sum_{j=0}^{\infty} U_j x^j = x F_U(x) \cdot \checkmark$$

Basics Definitions How to build Some visual exampler Structure Clustering Degree distributions Configuration model Largest component Generating Functions Definitions Properties

Random Networks

Frame 71/83 日 のへへ

Random Network

Basics

Generating

Properties

Connecting generating functions

- Goal: figure out forms of the component generating functions, *F_π* and *F_ρ*.
- π_n = probability that a random node belongs to a finite component of size n

$$= \sum_{k=0}^{\infty} P_k \times \Pr\left(\begin{array}{c} \text{sum of sizes of subcomponents} \\ \text{at end of } k \text{ random links} = n-1 \end{array}\right)$$

Therefore: $F_{\pi}(x) = \underbrace{x}_{SN2} \underbrace{F_{P}(F_{\rho}(x))}_{SN1}$

Extra factor of x accounts for random node itself.

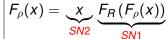
Frame 73/83

Connecting generating functions

- *ρ_n* = probability that a random link leads to a finite subcomponent of size *n*.
- Invoke one step of recursion: ρ_n = probability that a random node arrived along a random edge is part of a finite subcomponent of size n.

$$= \sum_{k=0}^{\infty} R_k \times \Pr\left(\begin{array}{c} \text{sum of sizes of subcomponents} \\ \text{at end of } k \text{ random links} = n-1 \end{array}\right)$$

Therefore:



 Again, extra factor of x accounts for random node itself.

Component sizes

Remembering vaguely what we are doing:

Finding F_P to obtain the size of the largest component $S_1 = 1 - F_{\pi}(1)$.

• Set x = 1 in our two equations:

 $F_{\pi}(1) = F_{P}(F_{\rho}(1))$ and $F_{\rho}(1) = F_{R}(F_{\rho}(1))$

- Solve second equation numerically for $F_{\rho}(1)$.
- Plug $F_{\rho}(1)$ into first equation to obtain $F_{\pi}(1)$.

Degree distributions Configuration model Largest component Generating Functions Definitions Proporties References

Frame 74/83

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Random Networks

Basics

Definitions

Structure

Generating

Frame 76/83

P

unctions

Properties

Random Networks

Basics

Structure

Connecting generating functions

We now have two functional equations connecting our generating functions:

 $F_{\pi}(x) = xF_{P}(F_{\rho}(x))$ and $F_{\rho}(x) = xF_{R}(F_{\rho}(x))$

- Taking stock: We know $F_P(x)$ and $F_R(x) = F'_P(x)/F'_P(1)$.
- We first untangle the second equation to find F_{ρ}
- We can do this because it only involves F_{ρ} and F_{R} .
- The first equation then immediately gives us F_π in terms of F_ρ and F_R.

Random Network

Basics

unctions

Properties

Component sizes

Example: Standard random graphs.

• We can show $F_P(x) = e^{-\langle k \rangle (1-x)}$

$$\therefore F_R(x) = F'_P(x)/F'_P(1) = e^{-\langle k \rangle (1-x)}/e^{-\langle k \rangle (1-x')}|_{x'=1}$$

$$= e^{-\langle k \rangle (1-x)} = F_P(x)$$
 ...aha!

RHS's of our two equations are the same.

• So
$$F_{\pi}(x) = F_{\rho}(x) = xF_R(F_{\rho}(x)) = xF_R(F_{\pi}(x))$$

Why our dirty (but wrong) trick worked earlier...

Random Networks

Basics

Structure

Generating

Properties

Frame 77/83

Component sizes

• We are down to $F_{\pi}(x) = xF_R(F_{\pi}(x))$ and $F_R(x) = xe^{-\langle k \rangle(1-x)}$.

$$\therefore$$
 $F_{\pi}(x) = x e^{-\langle k \rangle (1 - F_{\pi}(x))}$

We're first after S₁ = 1 − F_π(1) so set x = 1 and replace F_π(1) by 1 − S₁:

$$1-S_1=e^{-\langle k \rangle S_1}$$

- Just as we found with our dirty trick...
- Again, have to resort to numerics at this point.

Average component size

Example: Standard random graphs.

- Use fact that $F_P = F_R$ and $F_{\pi} = F_{\rho}$.
- Two differentiated equations reduce to only one:

$$F'_{\pi}(x) = F_{\mathcal{P}}\left(F_{\pi}(x)
ight) + xF'_{\pi}(x)F'_{\mathcal{P}}\left(F_{\pi}(x)
ight)$$

Rearrange:
$$F'_{\pi}(x) = \frac{F_P(F_{\pi}(x))}{1 - xF'_P(F_{\pi}(x))}$$

- Simplify denominator using $F'_{\pi}(x) = \langle k \rangle F_{\pi}(x)$
- Replace $F_P(F_\pi(x))$ using $F_\pi(x) = xF_P(F_\pi(x))$.
- Set x = 1 and replace $F_{\pi}(1)$ with $1 S_1$.

End result:
$$\langle n \rangle = F'_{\pi}(1) = \frac{(1-S_1)}{1-\langle k \rangle(1-S_1)}$$

Random Network
Basics Definitions How to build
Some visual examples Structure
Clustering Degree distributions Configuration model Largest component
Generating Functions Definitions Properties
References
Frame 78/83

日 りへへ

Random Networks

Basics

Properties

Frame 80/83

Average component size

- Next: find average size of finite components $\langle n \rangle$.
- Using standard G.F. result: $\langle n \rangle = F'_{\pi}(1)$.
- Try to avoid finding $F_{\pi}(x)$...
- Starting from $F_{\pi}(x) = xF_{P}(F_{\rho}(x))$, we differentiate:

 $F_{\pi}'(x)=F_{P}\left(F_{
ho}(x)
ight)+xF_{
ho}'(x)F_{P}'\left(F_{
ho}(x)
ight)$

• While $F_{\rho}(x) = xF_R(F_{\rho}(x))$ gives

$$F_{
ho}'(x)=F_R(F_{
ho}(x))+xF_{
ho}'(x)F_R'(F_{
ho}(x))$$

- Now set x = 1 in both equations.
- We solve the second equation for F'_ρ(1) (we must already have F_ρ(1)).
- Plug $F'_{\rho}(1)$ and $F_{\rho}(1)$ into first equation to find $F'_{\pi}(1)$.

Average component size

Our result for standard random networks:

$$\langle n \rangle = F'_{\pi}(1) = \frac{(1-S_1)}{1-\langle k \rangle(1-S_1)}$$

- Recall that (k) = 1 is the critical value of average degree for standard random networks.
- Look at what happens when we increase (k) to 1 from below.
- We have $S_1 = 0$ for all $\langle k \rangle < 1$ so

$$\langle n \rangle = \frac{1}{1 - \langle k \rangle}$$

- This blows up as $\langle k \rangle \rightarrow 1$.
- Reason: we have a power law distribution of component sizes at (k) = 1.
- Typical critical point behavior....

Frame 81/83

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Some visual example Structure Clustering Degree distributions Configuration model Largest component Generating Functions Definitions Properties

Random Network

Basics

Random Network

Basics

Properties

Average component size

• Limits of $\langle k \rangle = 0$ and ∞ make sense for

$$\langle n
angle = F'_{\pi}(1) = rac{(1-S_1)}{1-\langle k
angle(1-S_1)}$$

- As $\langle k \rangle \rightarrow 0$, $S_1 = 0$, and $\langle n \rangle \rightarrow 1$.
- All nodes are isolated.
- As $\langle k \rangle \to \infty$, $S_1 \to 1$ and $\langle n \rangle \to 0$.
- ▶ No nodes are outside of the giant component.



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Random Networks

References

Frame 83/83 日 のへへ